

# Network Alignment by Discrete Ollivier-Ricci Flow

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**Speaker: Kin Sum Liu<sup>3</sup>**

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## An Example

Consider the same group of people who participate in two social network platforms:

- Private network: identities not revealed, e.g., Facebook .

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Assume these two networks have **almost the same** topology.

Goal: align the two networks by vertex correspondences, hence reveals the identities of the private network.

# Graph Isomorphism

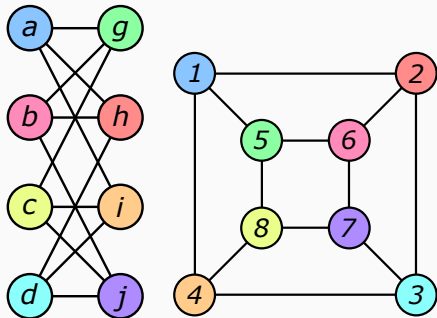
Given a pair of graphs  $G_1, G_2$ , find a one-to-one correspondence of the vertices in  $G_1$  to vertices in  $G_2$  such that  $(u, v)$  is an edge in  $G_1$  if and only if their corresponding nodes  $f(u), f(v)$  are connected in  $G_2$ <sup>1</sup>.

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- Subgraph isomorphism is NP-complete.

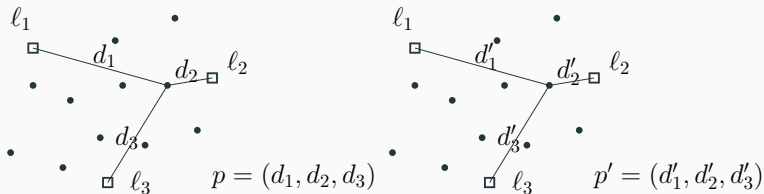
# Graph Isomorphism

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- Many practical algorithms: e.g., NAUTY.
- Subgraph isomorphism is NP-complete.
- **Approximate graph isomorphism**: find the best correspondence between vertices in  $G_1$  and  $G_2$  s.t. if  $u, v$  are connected in  $G_1$  their corresponding nodes are likely connected in  $G_2$ .

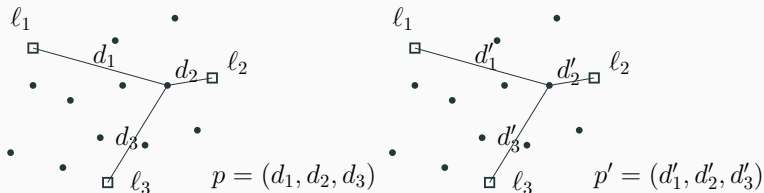
# Our Solution: A Geometric Embedding Approach

How to align two sets of points in some embedding plane, assuming that some landmarks  $\ell_i$  are already aligned?



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- Any point  $p$  can be represented by the barycentric coordinates  $(d_1, d_2, d_3)$ ,  $d_i$  is distance to  $\ell_i$ .
- If the barycentric coordinates of  $p$  and  $p'$  are similar, we match  $p$  and  $p'$ .

## Quantify the 'Position' of a Node in a Network

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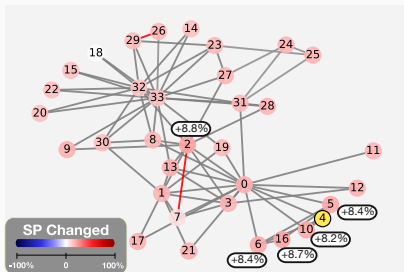
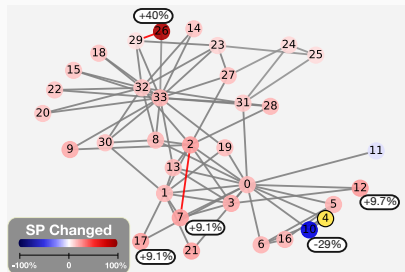
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Q: Robust to noises (edge insertion/deletion)?

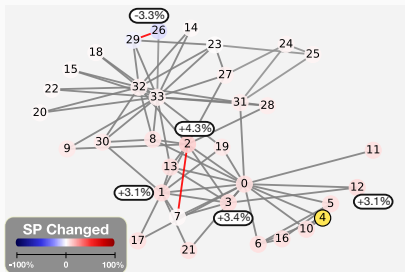
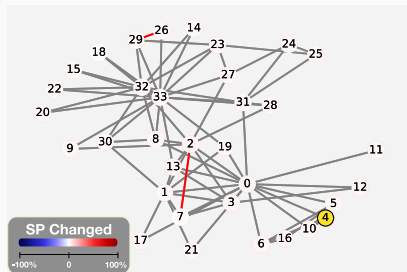
# Robustness: Remove Two Edges

Left: Spectral embedding; Right: Tutte/Spring embedding.



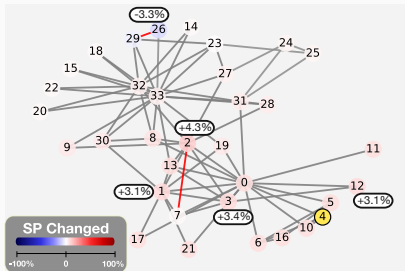
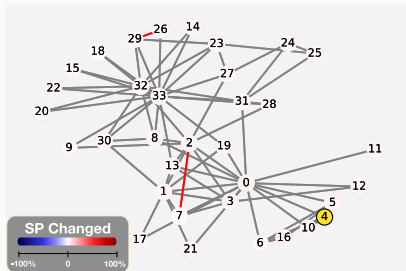
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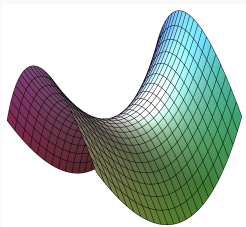
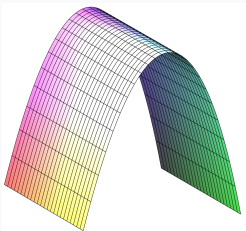
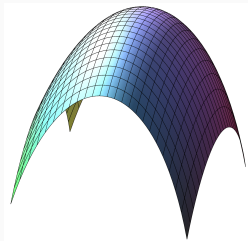
Q: How is our metric defined?

# Discrete Ricci Curvature & Ricci Flow



# Curvature in Geometry

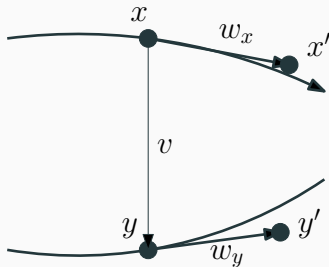
- Sphere: positive curvature;
- Plane: zero curvature;
- Hyperbolic plane: negative curvature.



# Sectional Curvature

Consider a tangent vector  $v = xy$ . Take another tangent vector  $w_x$  and transport it along  $v$  to be a tangent vector  $w_y$  at  $y$ .

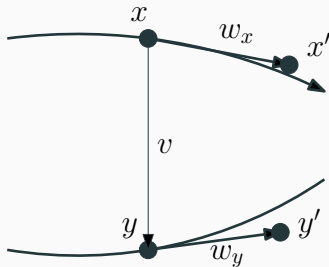
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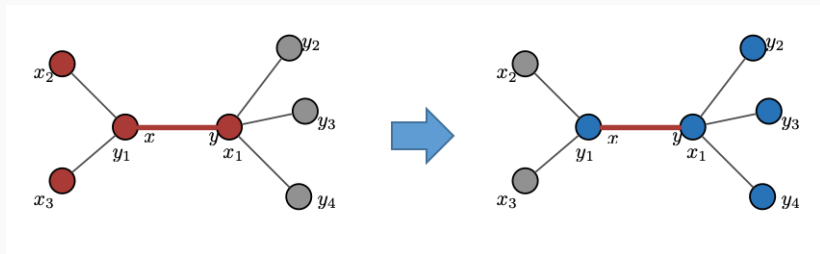
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- Ricci Curvature: averaging over all direction  $w$ .

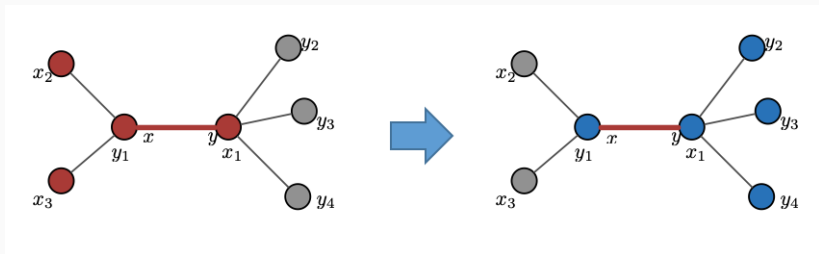
# Discrete Ricci Curvature

Take the analog: for an edge  $xy$ , consider the distances from  $x$ 's **neighbors** to  $y$ 's **neighbors** and compare it with the length of  $xy$ .



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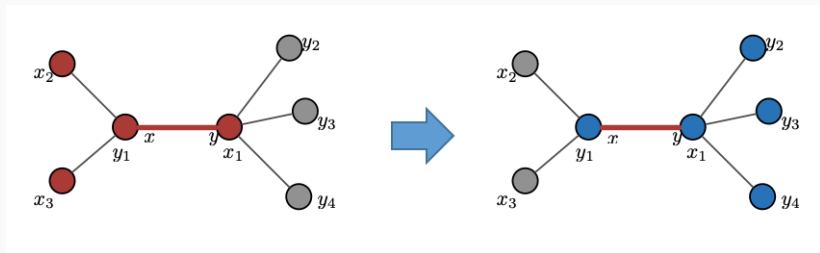
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# Discrete Ricci Curvature

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- Issue: how to match  $x$ 's neighbors to  $y$ 's neighbors?
- Assign uniform distribution  $\mu_1, \mu_2$  on  $x$ ' and  $y$ 's neighbors.
- Use optimal transportation distance (earth-mover distance) from  $\mu_1$  to  $\mu_2$ : the matching that minimize the total transport distance.

## Definition (Ollivier)

Let  $(X, d)$  be a metric space and let  $m_1, m_2$  be two probability measures on  $X$ . For any two distinct points  $x, y \in X$ , the (Ollivier-) Ricci curvature along  $xy$  is defined as

$$\kappa(x, y) := 1 - \frac{W_1(m_x, m_y)}{d(x, y)},$$

where  $m_x$  ( $m_y$ ) is a probability distribution defined on  $x$  ( $y$ ) and its neighbors,  $W_1(\mu_1, \mu_2)$  is the  $L_1$  **optimal transportation distance** between two probability measure  $\mu_1$  and  $\mu_2$  on  $X$ :

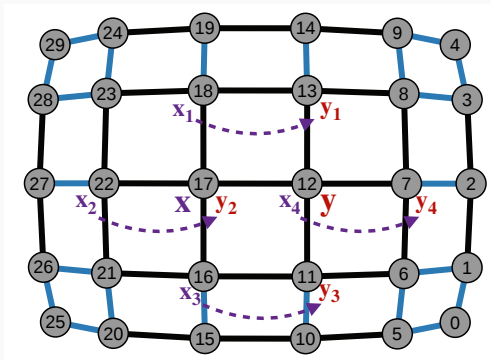
$$W_1(\mu_1, \mu_2) := \inf_{\psi \in \Pi(\mu_1, \mu_2)} \int_{(u, v)} d(u, v) d\psi(u, v)$$

For a node  $w$  with  $k$  neighbors, we define

$m_w(w) = \alpha$ ;  $m_w(v) = (1 - \alpha)/k$ . We choose  $\alpha = 1/2$ .

# Examples

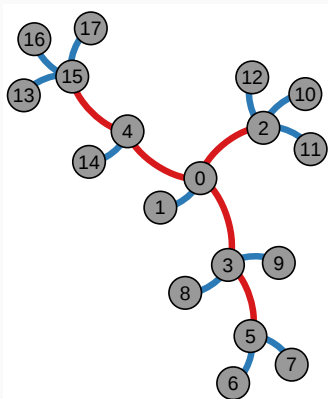
Zero curvature: 2D grid.





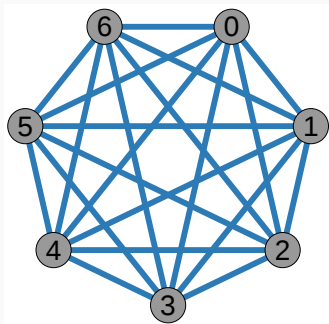
## Examples

Negative curvature: tree:  $\kappa(x, y) = 1/d_x + 1/d_y - 1$ ,  $d_x$  is degree of  $x$ .



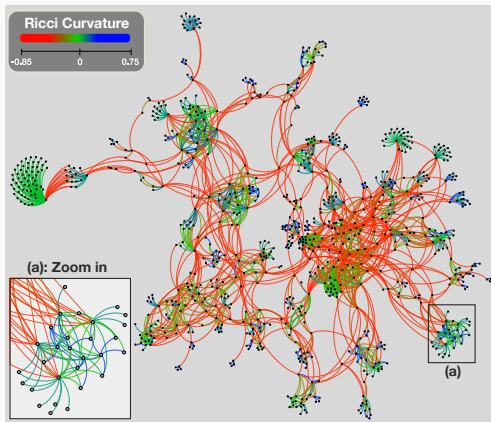
## Examples

Positive curvature: complete graph.



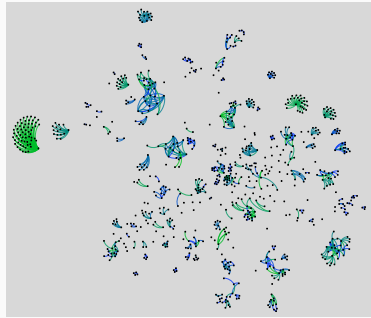
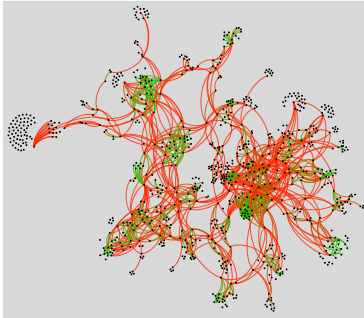
## Example: Ricci Curvature

Negatively curved edges are like “backbones”, maintaining the connectivity of clusters, in which edges are mostly positively curved.



# Curvature Distribution

Left: Negative curvature edges. Right: Positive curvature edges.<sup>2</sup>



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<sup>2</sup>ForceAtlas layout by Gephi

## Edge Weights Generated by Ricci flow

Given a graph  $G$  in which  $d(x, y)$  is the weight of the edge  $xy$  and  $\kappa(x, y)$  is the discrete Ricci curvature, we run

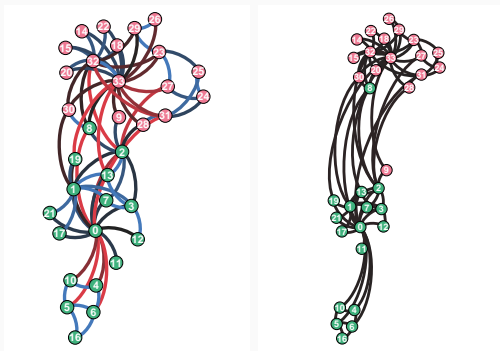
$$d_{i+1}(x, y) = (d_i(x, y) - \varepsilon \cdot \kappa_i(x, y) \cdot d_i(x, y)) \cdot N$$

Until convergence, where  $N$  is to rescale to make sure total edge weights remain the same.

At the limit,  $W(x, y)/d(x, y)$  is the same for all edges.

# Ricci Flow Metric

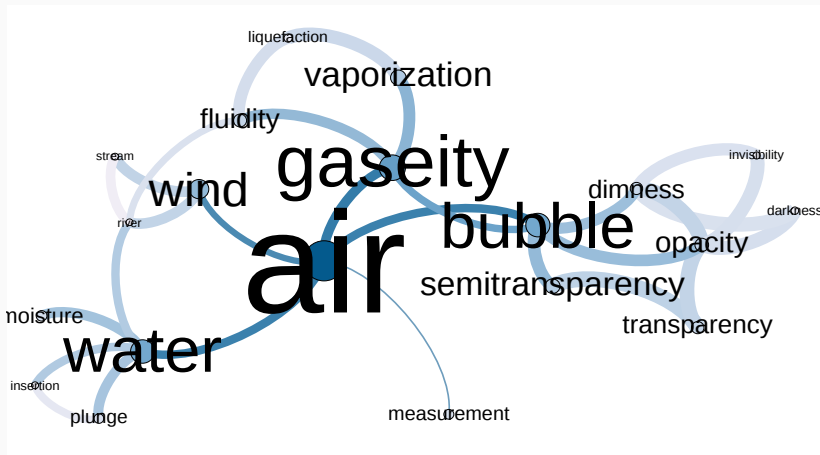
Intuition: flatten the network – shrink an edge if it is within a well connected community; stretch an edge if otherwise, s.t., the network curvature is uniform everywhere.<sup>3</sup>



<sup>3</sup>Karate Club by Gephi ForceAtlas layout

# Ricci Flow Metric on Semantic Wordnet

**As similarity metric:** On wordnet, edges between similar words are shrank s.t. similar words are closer with Ricci Flow Metric.



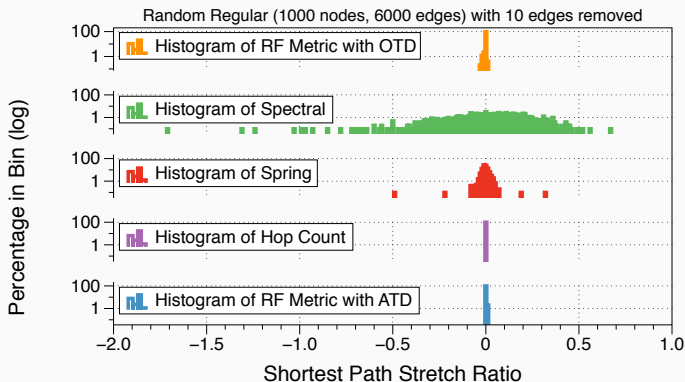
**Table 1:** Node similarity: Word distance by RF-Metric and hop count

Word	RF-Metric	Hop		Word	RF-Metric	Hop
<b>air</b>	0	0		<b>heaven</b>	0	0
gaseity	1.084512	1		hell	0.476738	1
bubble	1.233986	1		pleasure	0.673406	1
water	1.241377	1		pleasurableness	0.786310	1
wind	1.560098	1		hope	0.920200	1
vaporization	1.854184	2		pain	1.104767	2
semitransparency	1.900589	2		cheerfulness	1.253568	2
opacity	1.993095	2		content	1.254039	2
fluidity	2.032685	2		restoration	1.391618	1
transparency	2.077700	3		sweetness	1.432170	2
dimness	2.084738	2		physical pleasure	1.450673	2
moisture	2.204766	2		feeling	1.471766	2



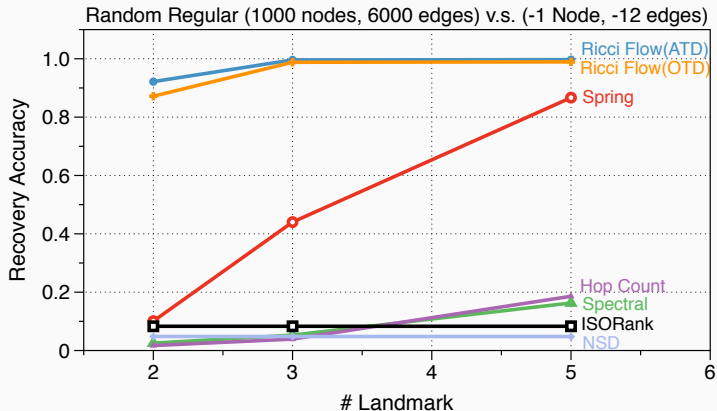
# Evaluation on Resilience

Randomly remove 10 edges in a random regular graph.



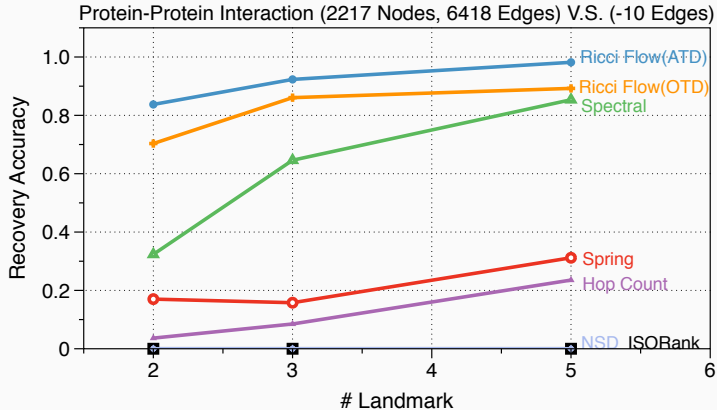
# Evaluation on Matching Performance

- Randomly remove one node in a random regular graph w/ degree 12.



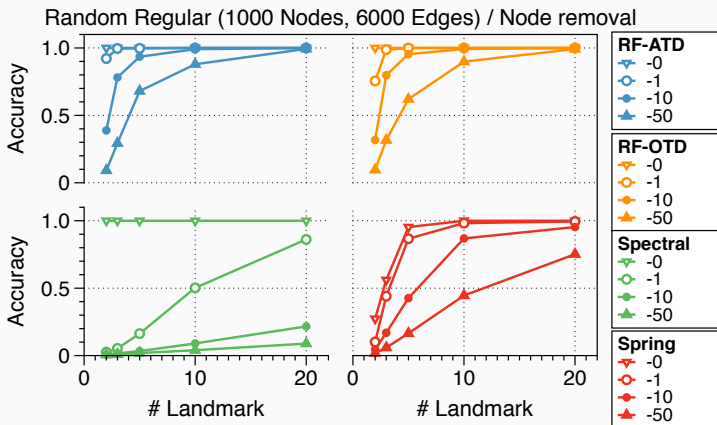
# Evaluation on Matching Performance

- Randomly remove 10 edges in a protein protein network.



# Evaluation on Matching Performance

- Random Regular Graph - remove Nodes



# Conclusions

Ricci flow metric on graph:

- A geometric metric that is robust to noises.
- Only require topology information to compute.
- Highly related to node similarity.

Ricci Curvature & Ricci Flow Source code Available:

<https://github.com/saibalmars/GraphRicciCurvature>

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