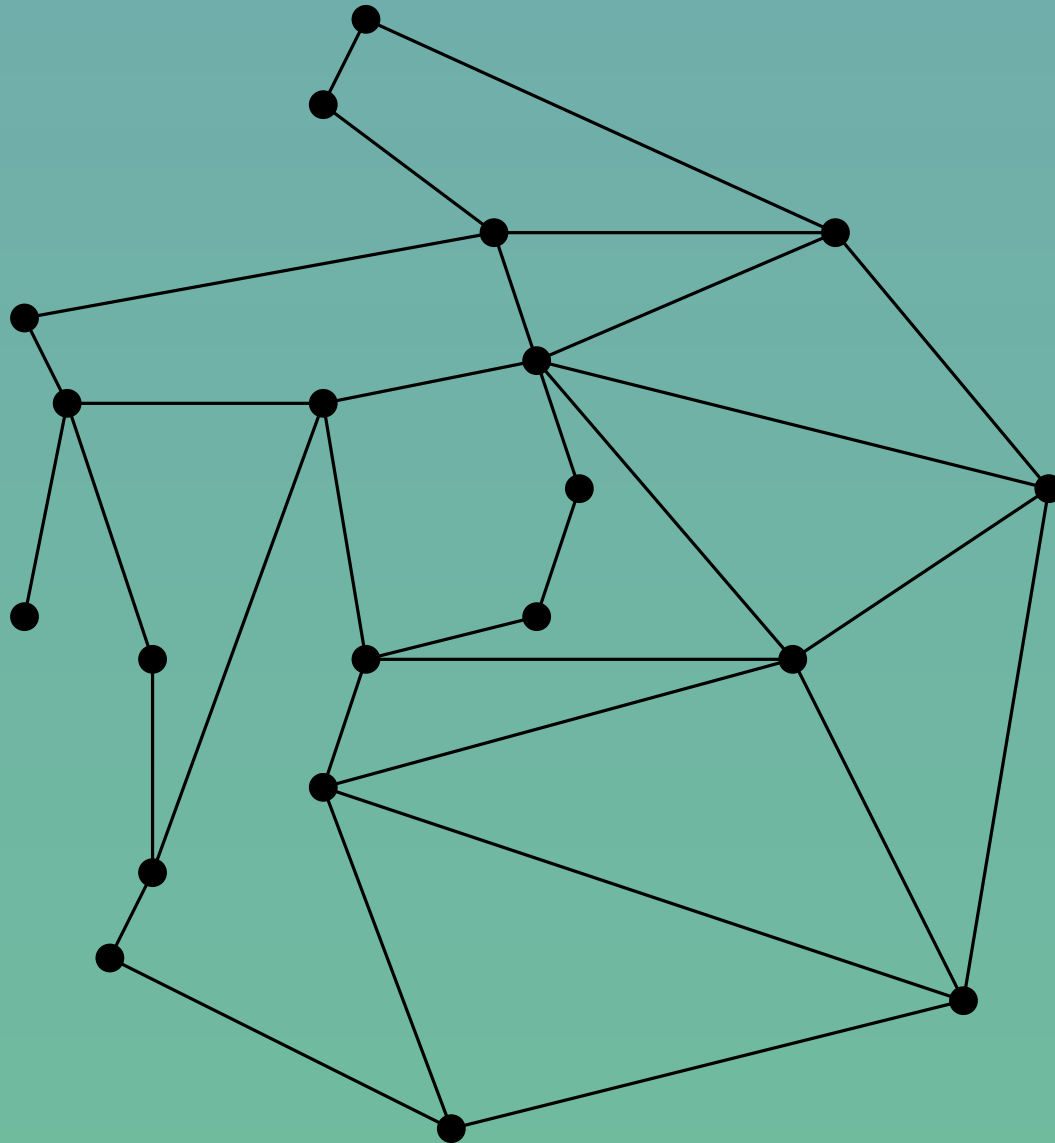


$\beta$ -Stars  
or  
On Extending a Drawing  
of a Connected Subgraph

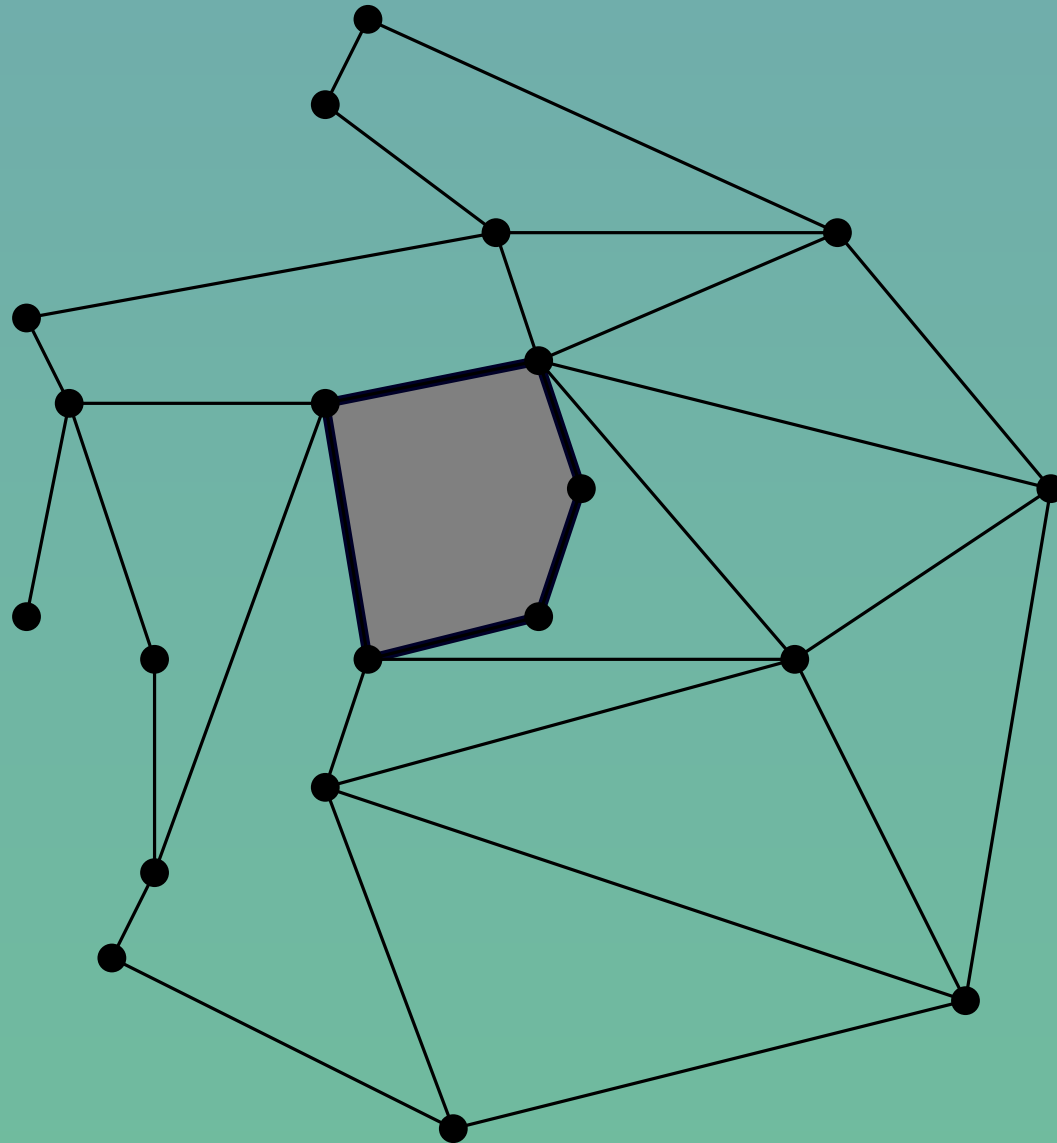
Tamara Mchedlidze  
Karlsruhe Institute of Technology

**Jérôme Urhausen**  
Utrecht University

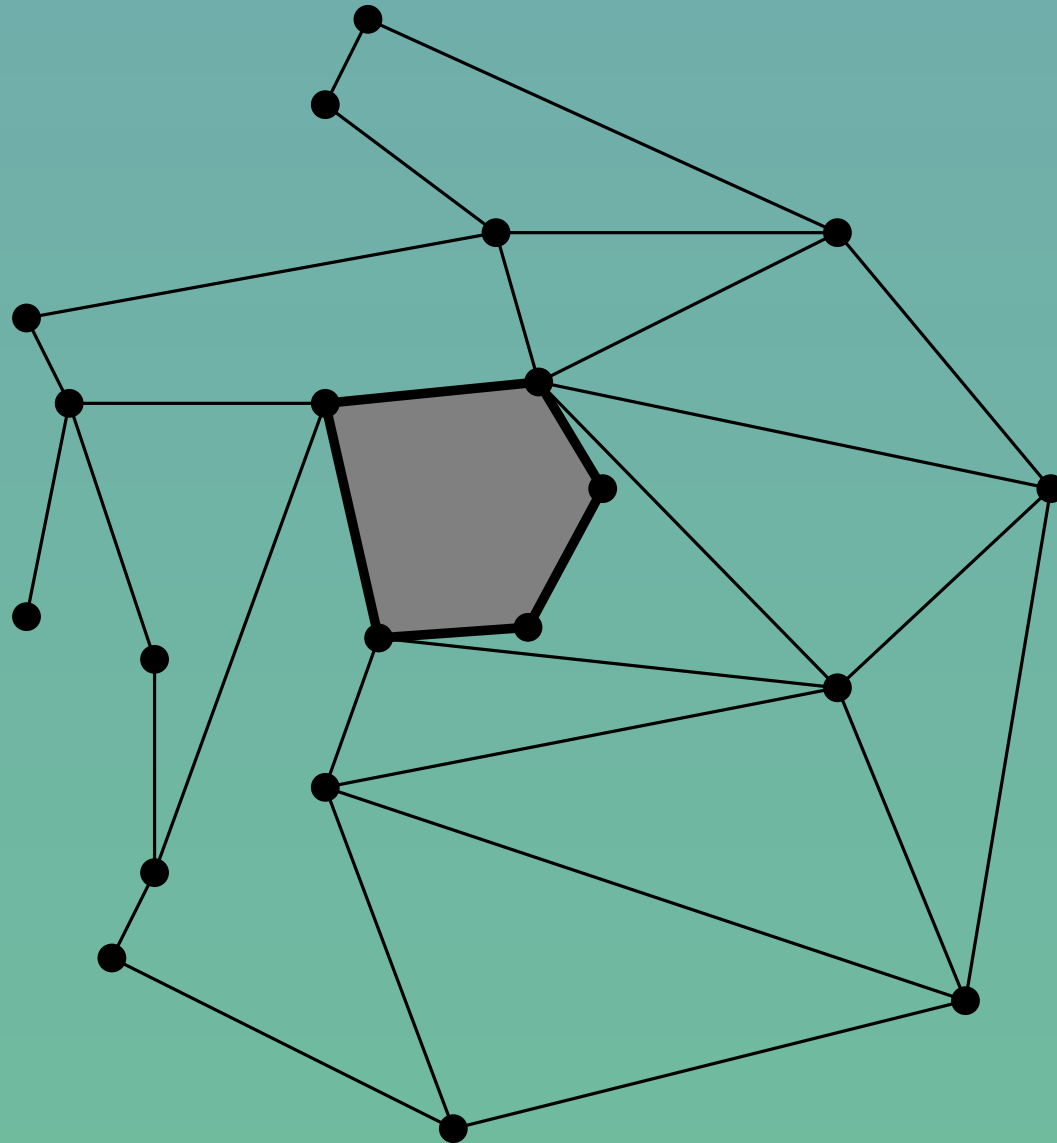
# The Problem



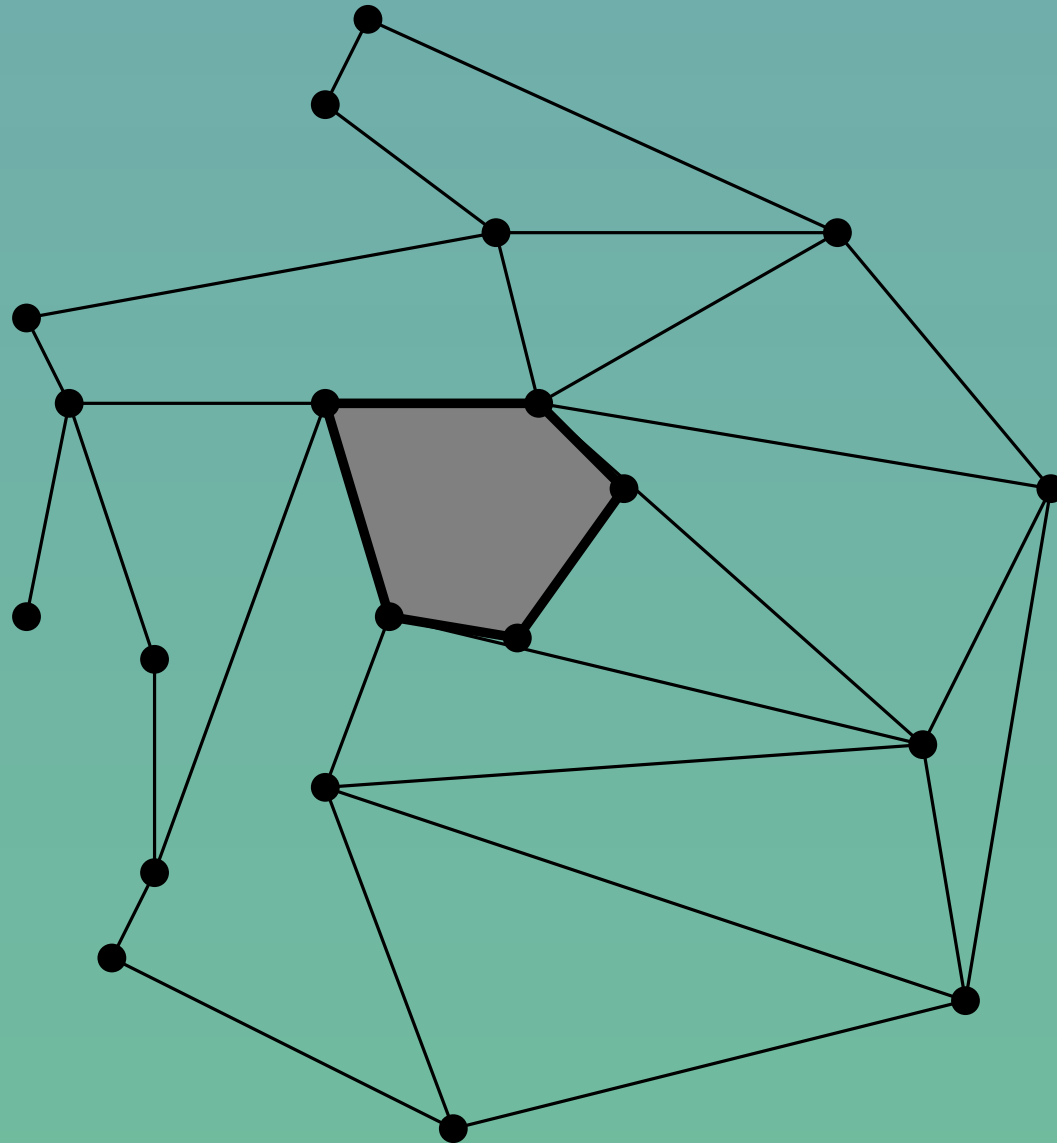
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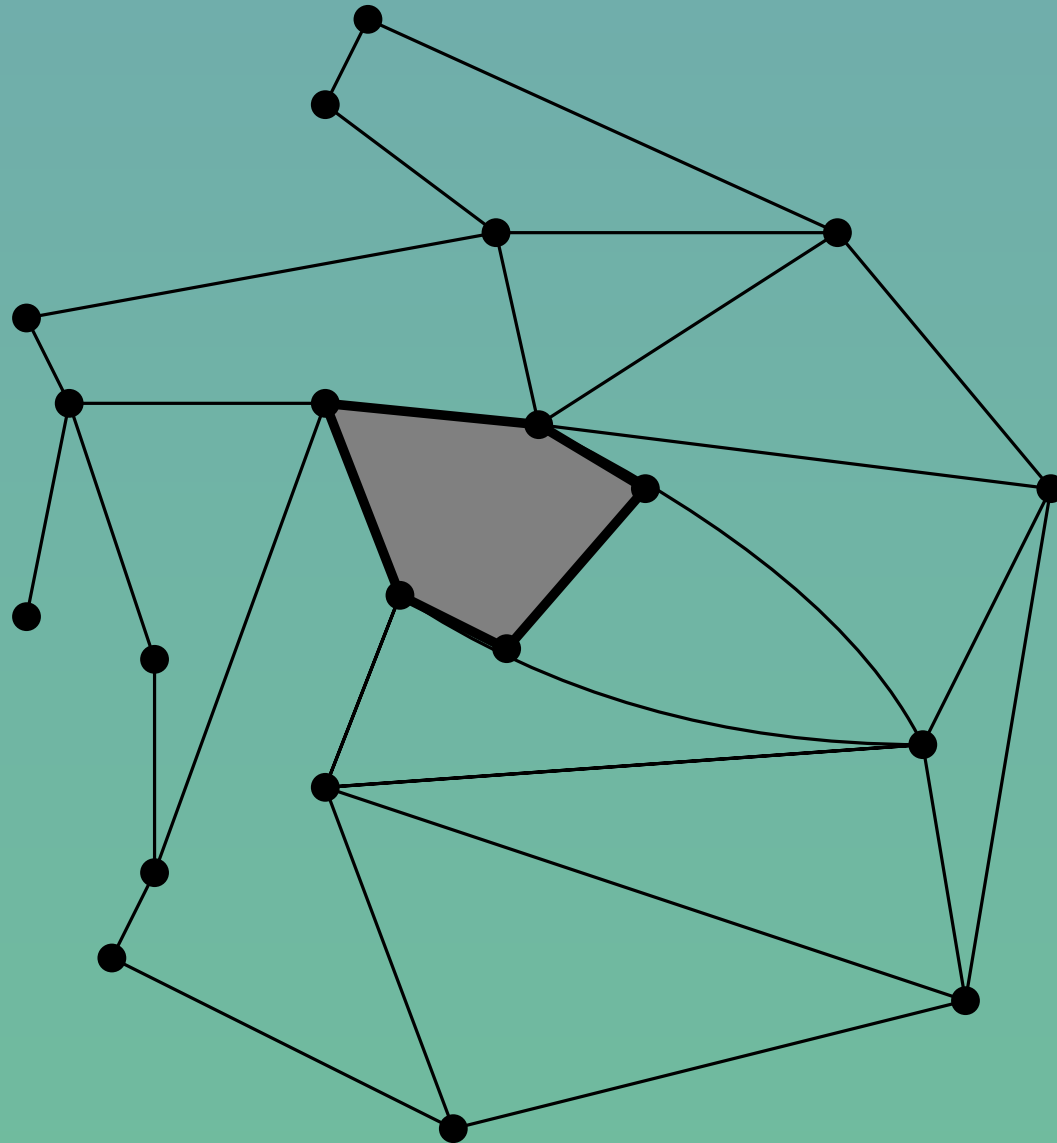
# The Problem



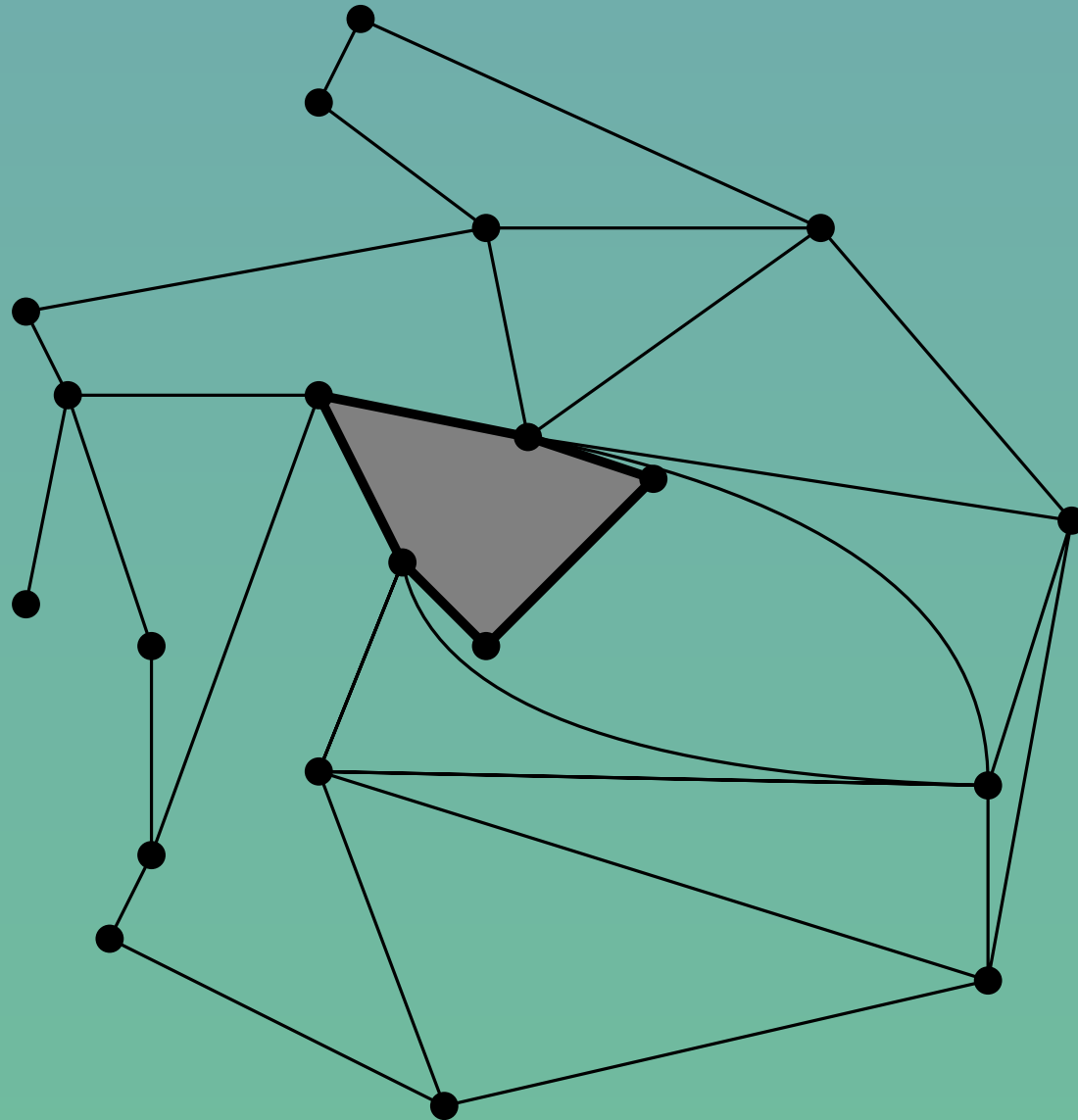
# The Problem



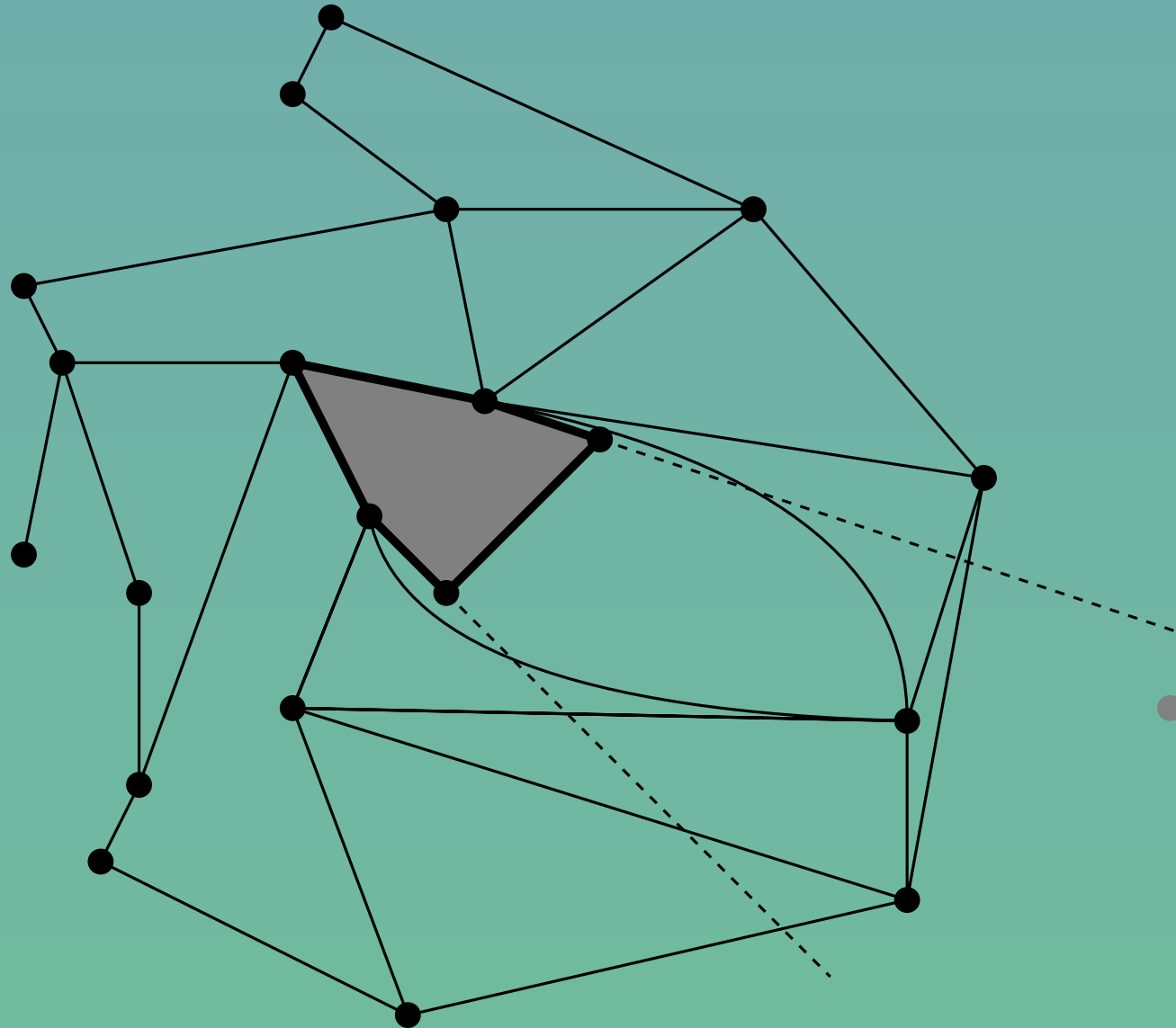
# The Problem



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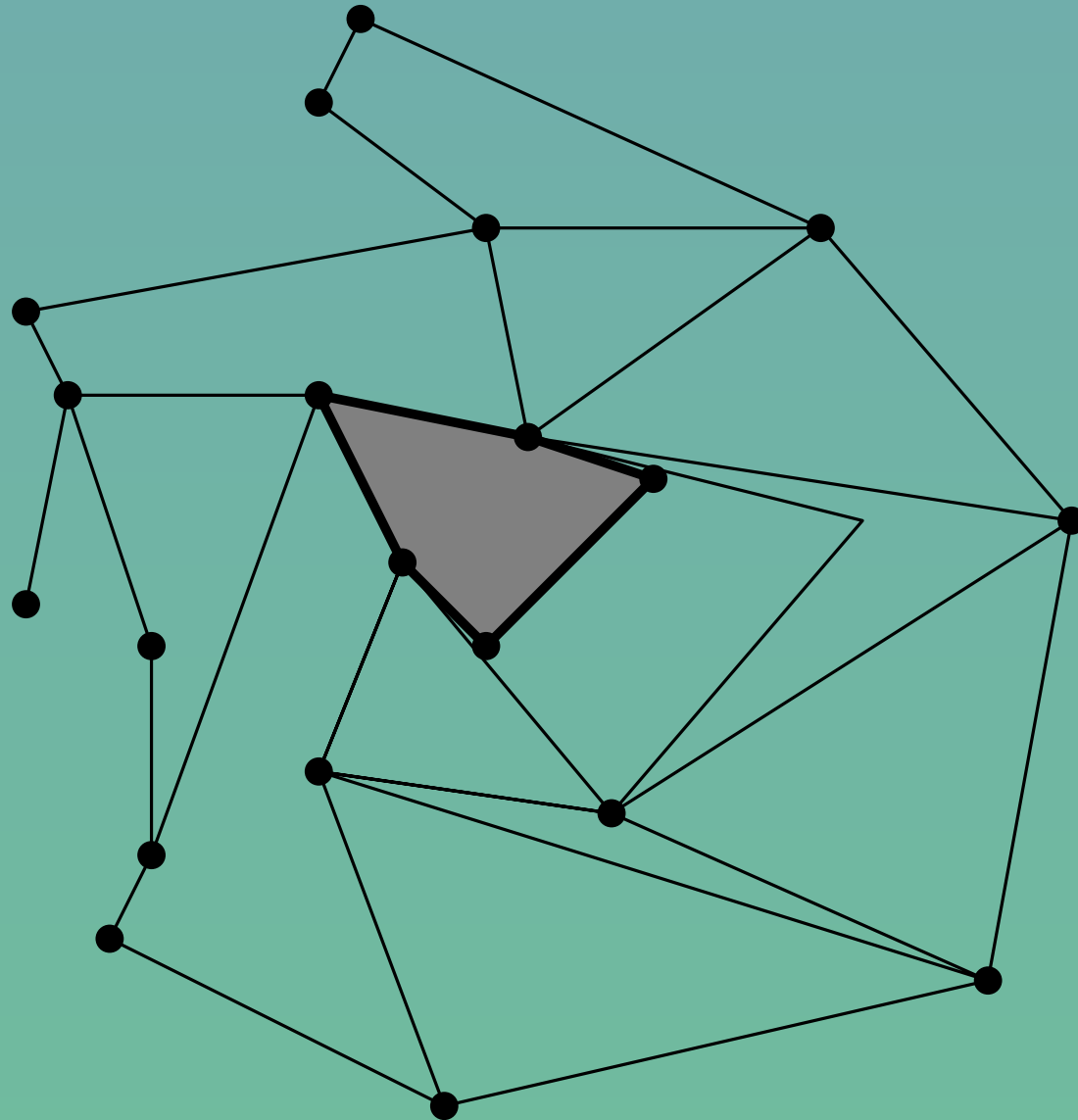


# The Problem





# The Problem



## Results

Given a plane graph  $G = (V, E)$  and a drawing of a face  $H$ , can the drawing be extended using straight lines?

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$H$  is outer face

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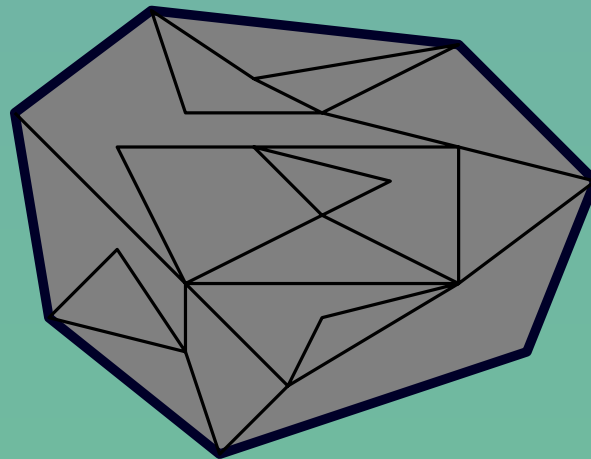
---

convex

yes [T63], [CEGL12]

star-shaped

---



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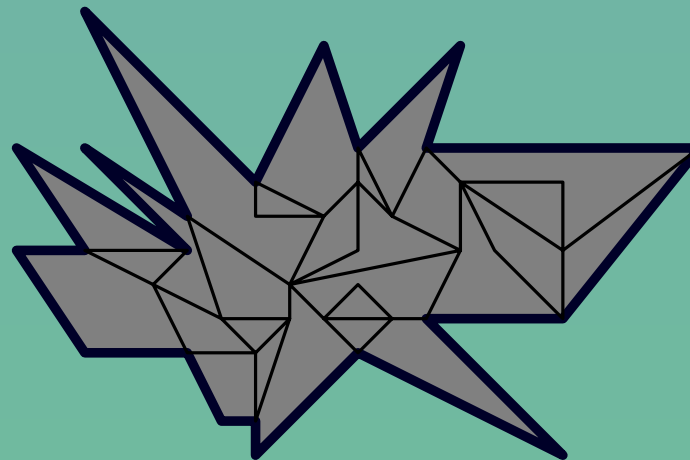
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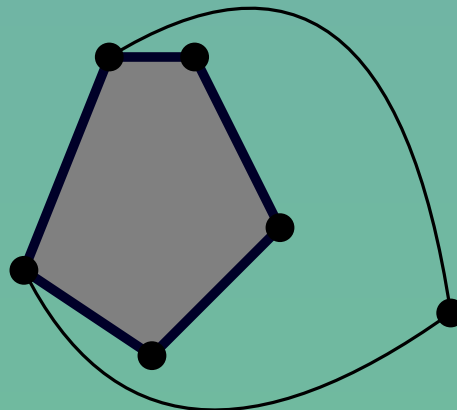
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	$H$ is outer face	$H$ is inner face
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star-shaped	yes [HN08]	indication of complexity



# Results

Given a plane graph  $G = (V, E)$ , a drawing of a subgraph  $H$ , extend the drawing using as few bends per edge as possible.

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# of bends per edge

---

$H = (V, \emptyset)$

$120|V|$  [PW01]

---



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$$H = (V, \emptyset) \quad 3|V| + 2 \text{ [BGL08]}$$

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-

$$72|H| \text{ [CFGLMS15]}$$

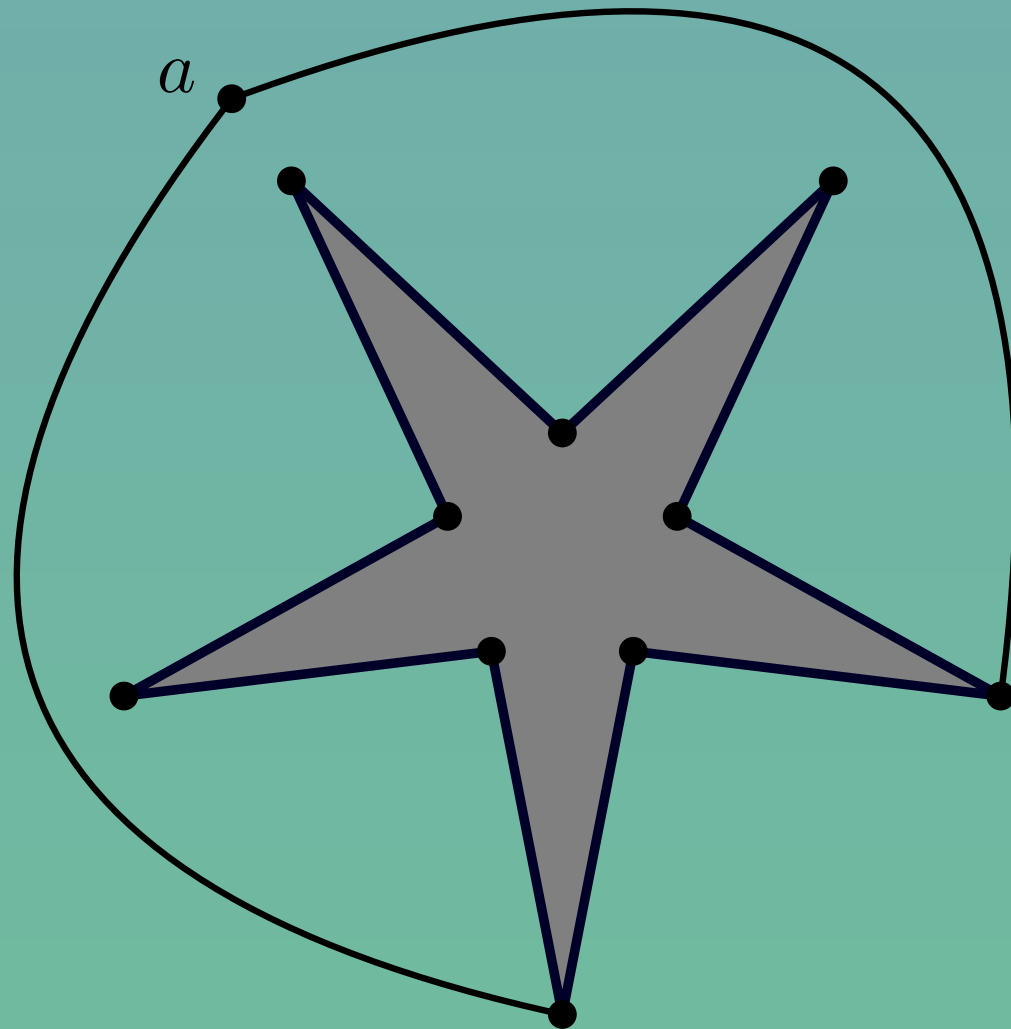
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# Results

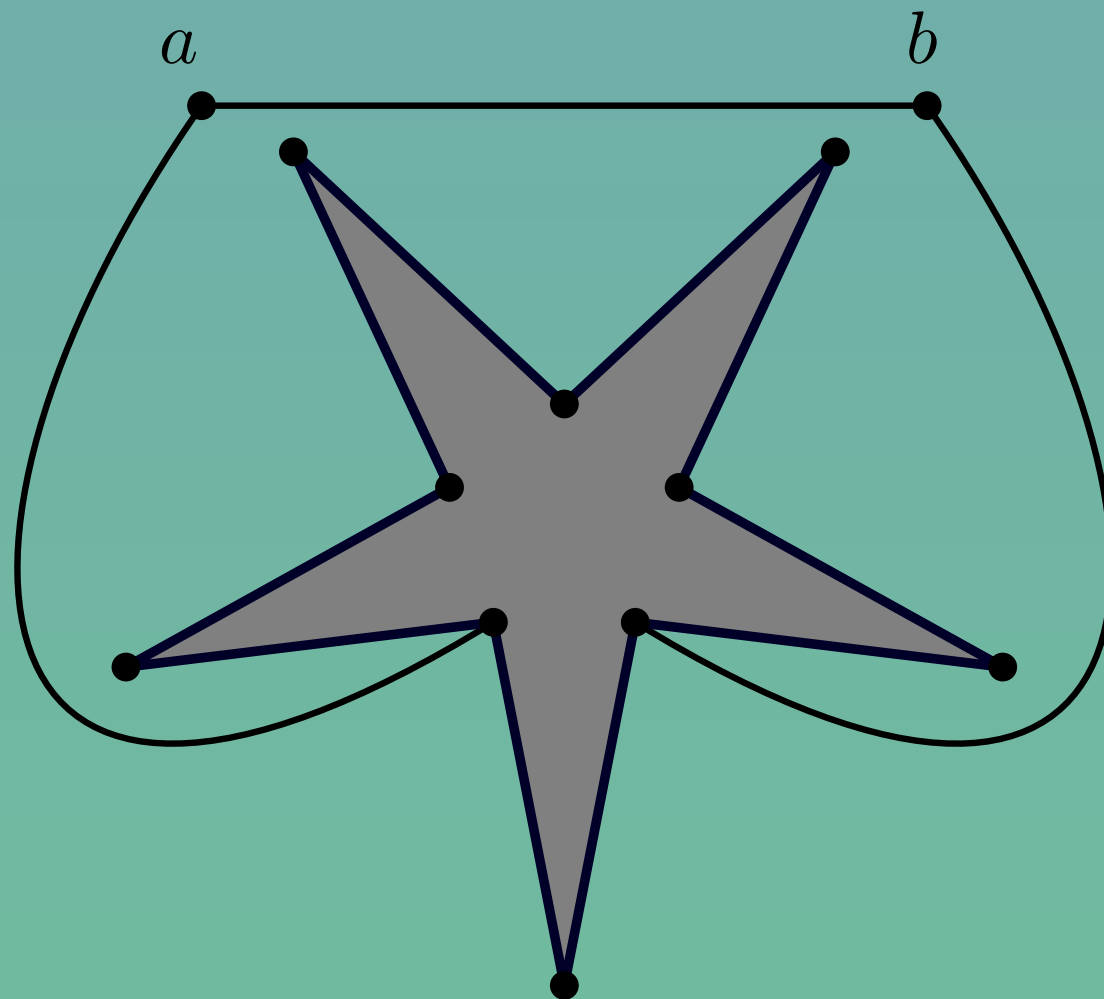
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	# of bends per edge
$H = (V, \emptyset)$	$3 V  + 2$ [BGL08]
-	$72 H $ [CFGLMS15]
$H$ connected	tight algorithm

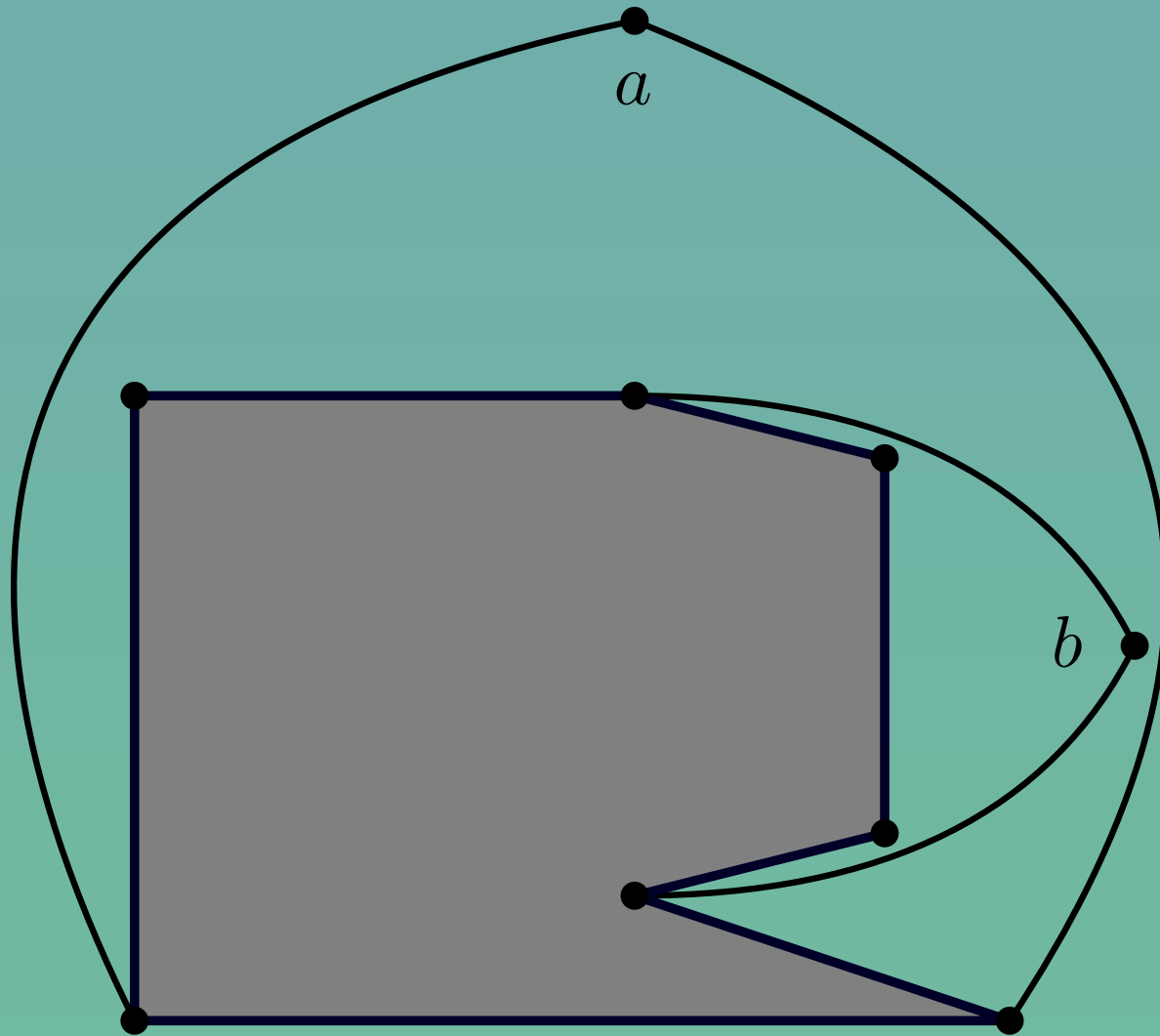
# Extending Stars



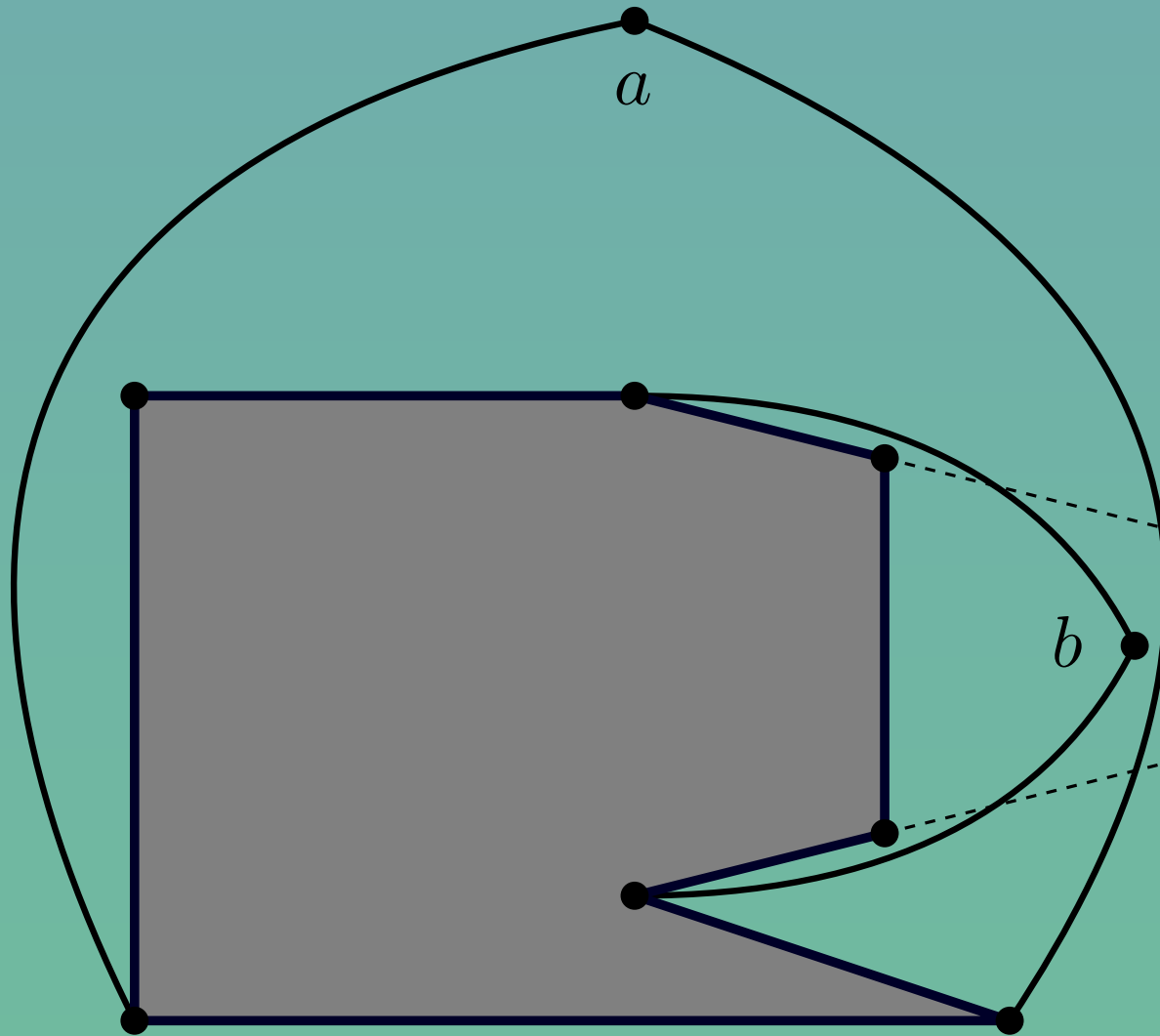
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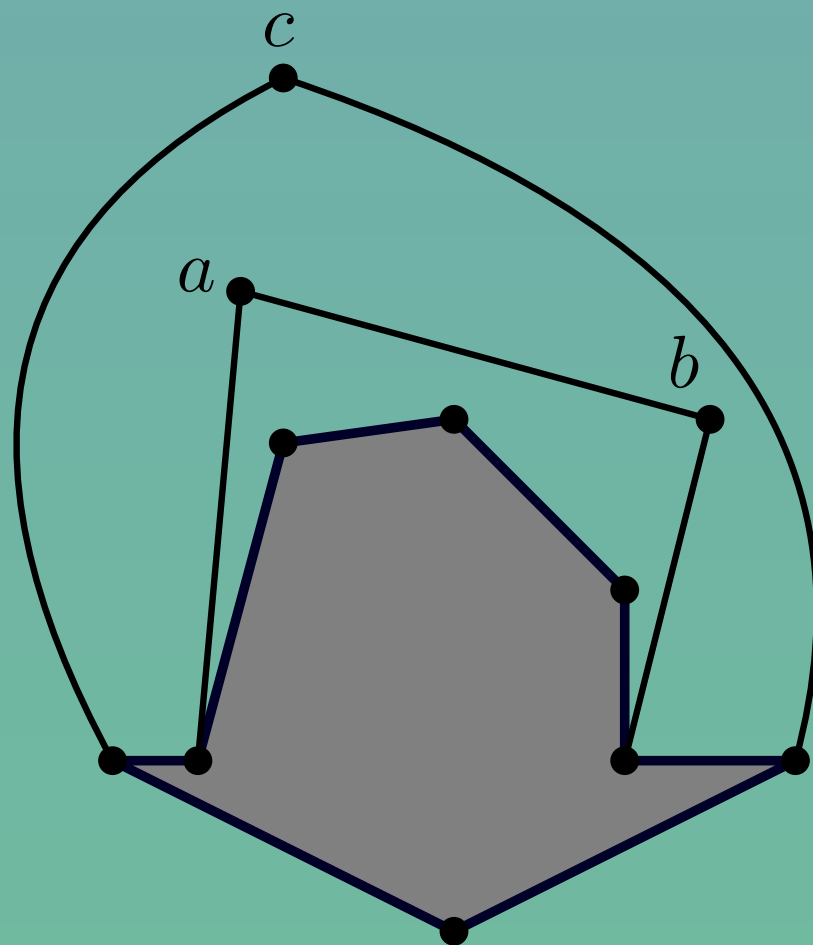
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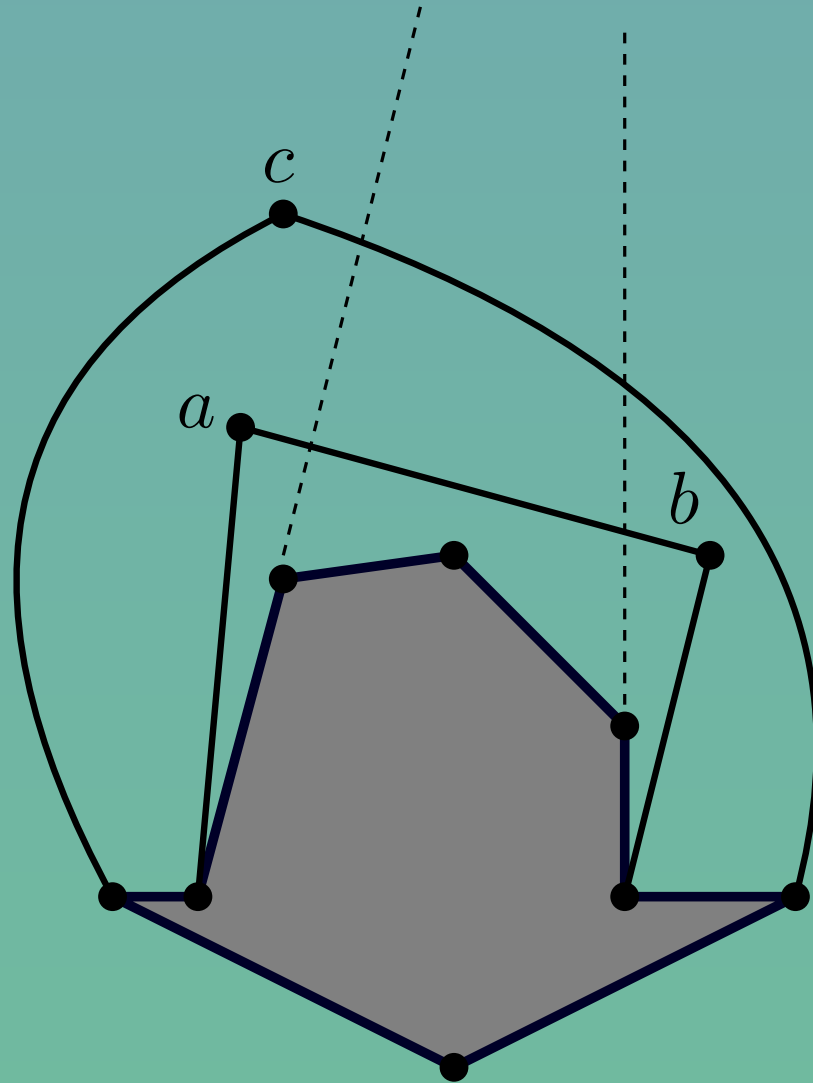


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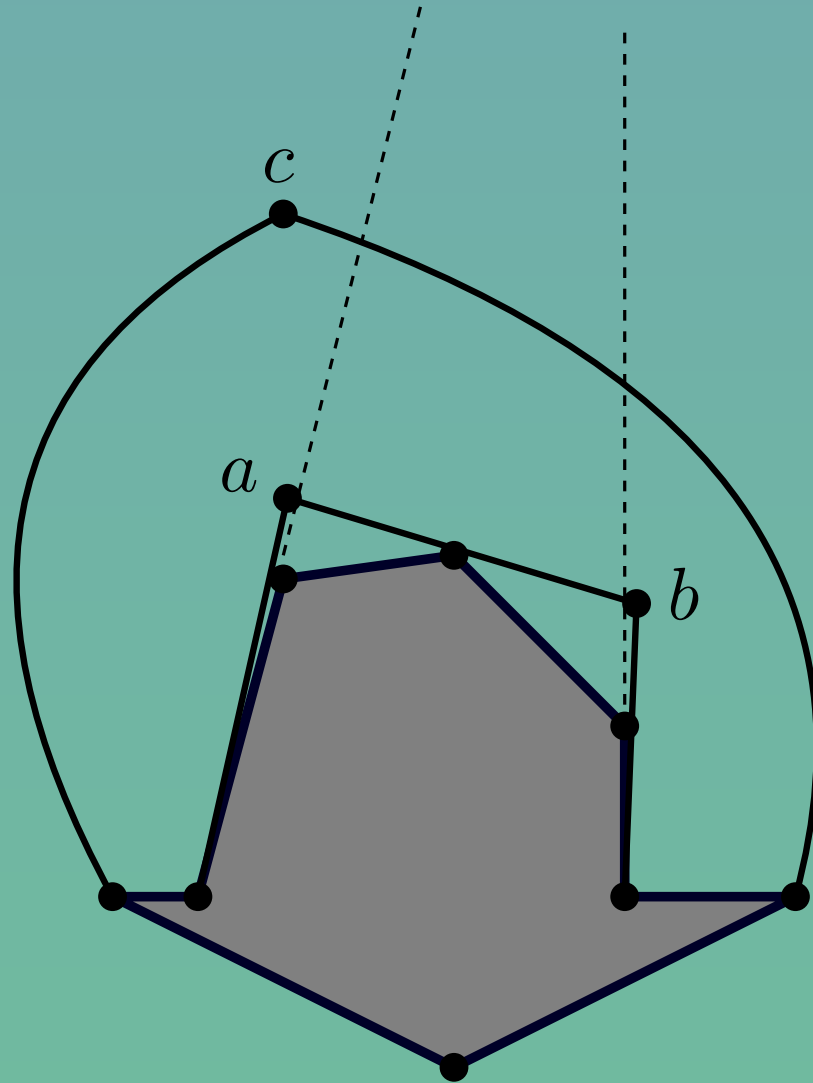




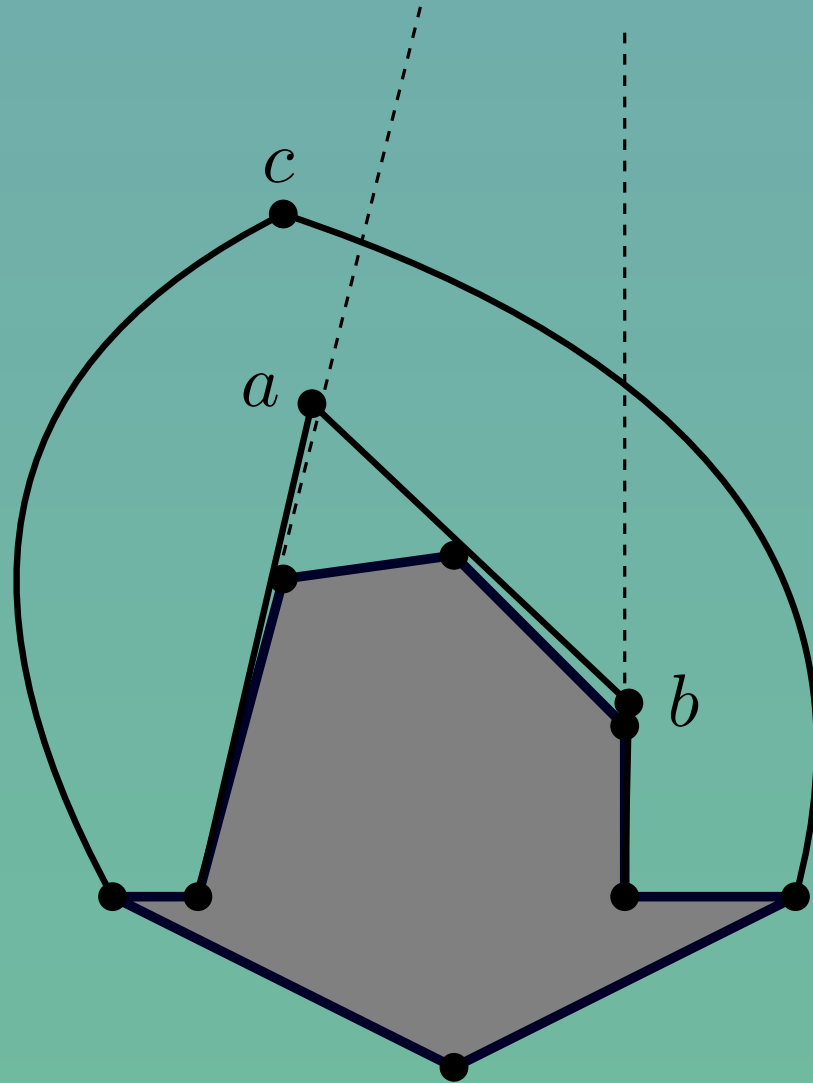
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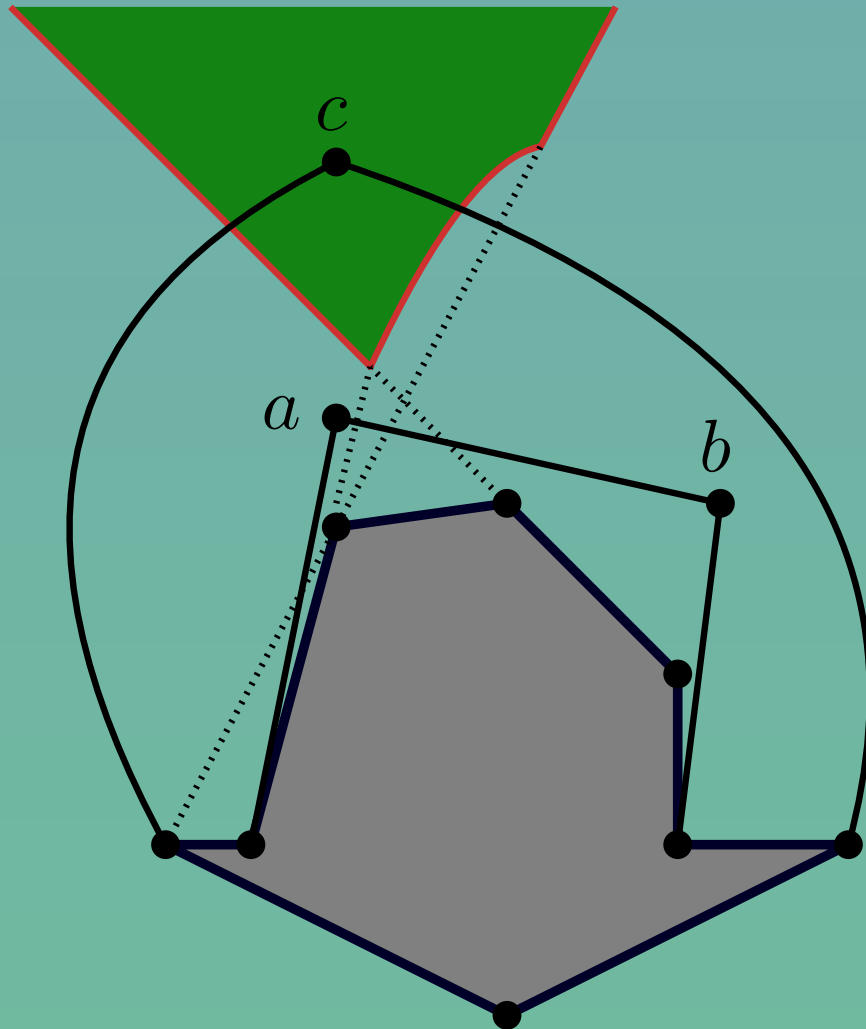
# Extending Stars



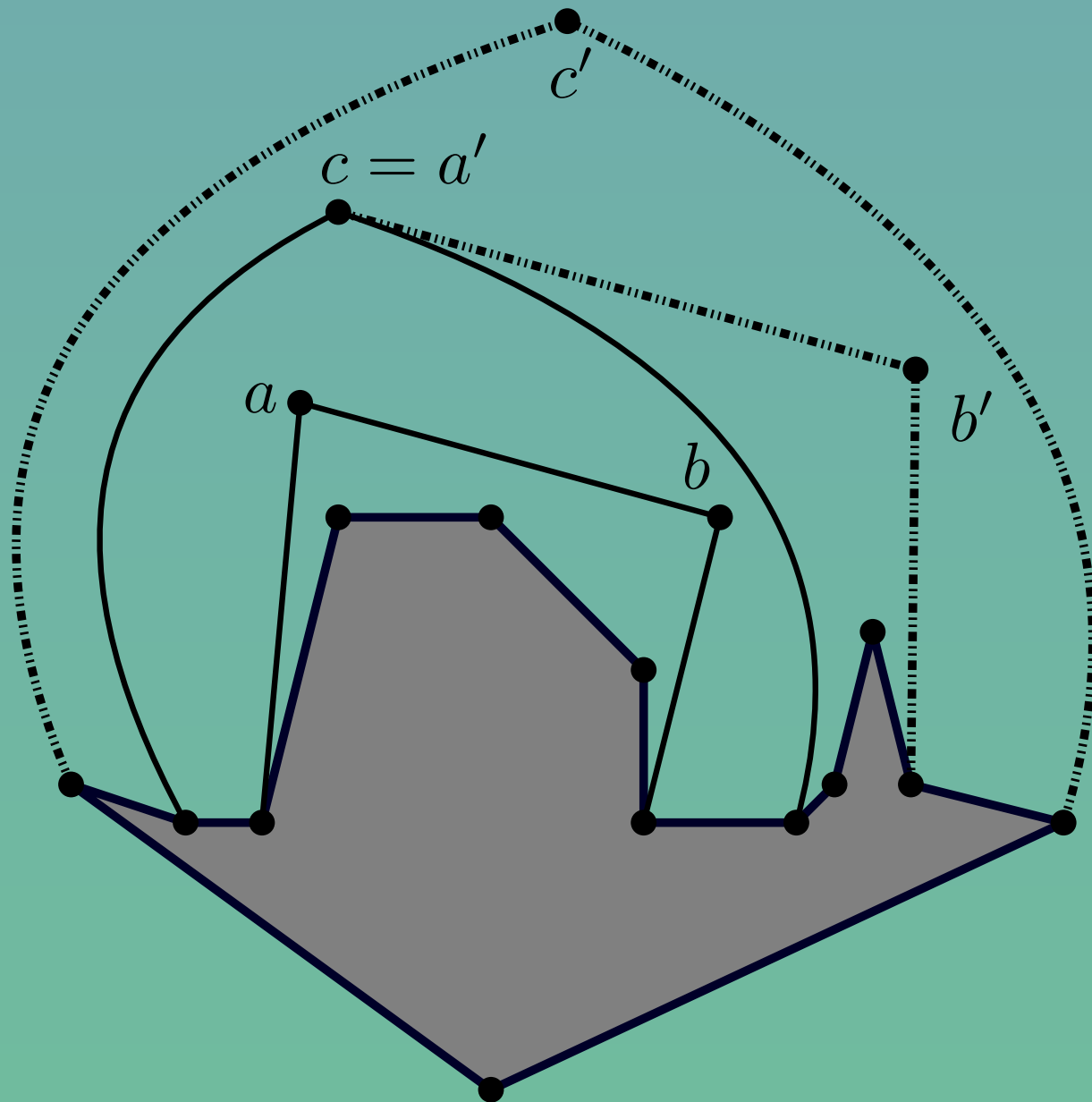
# Extending Stars



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# Extending Stars

## Theorem

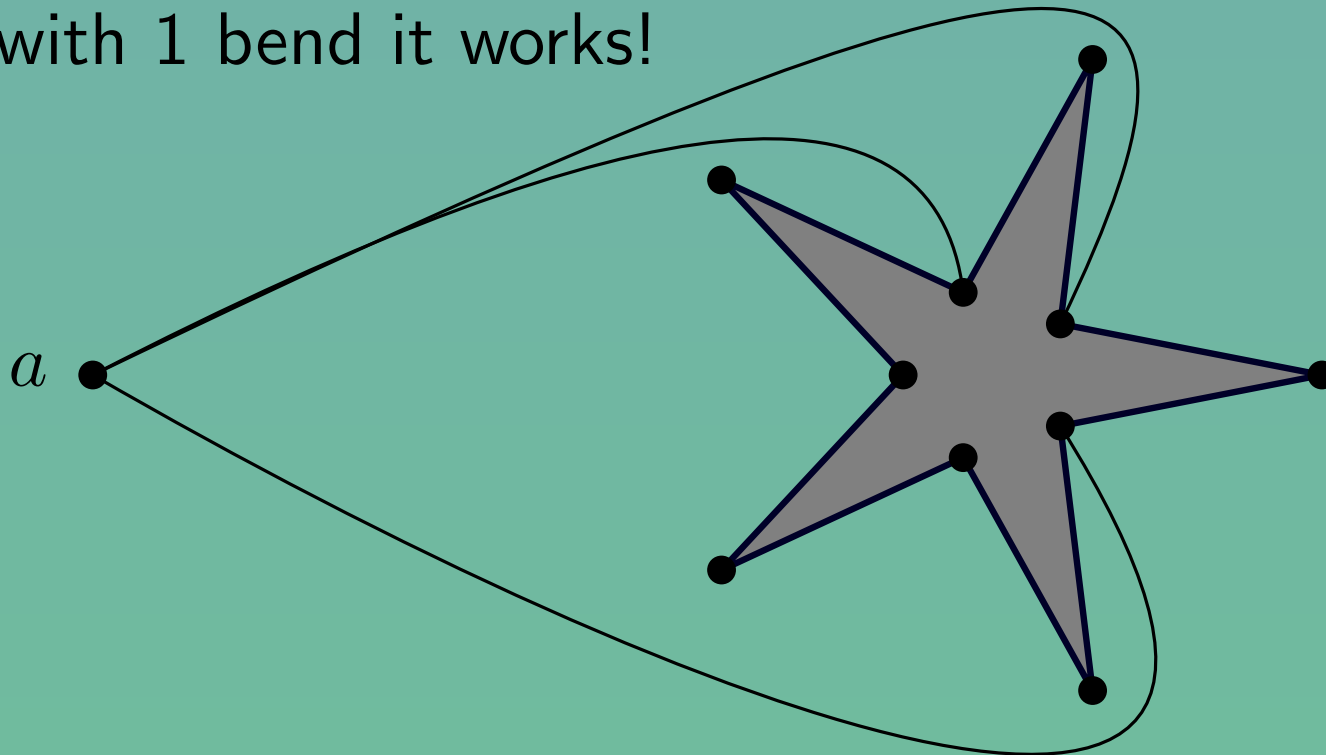
There is an instance  $(G, \Gamma_H)$  where  $\Gamma_H$  is a star-shaped inner face, such that the feasibility area of some vertex  $v \in G$  is partially bounded by a curve whose implicit representation is a polynomial of degree  $2^{\Omega(|V|)}$ .

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**But** with 1 bend it works!

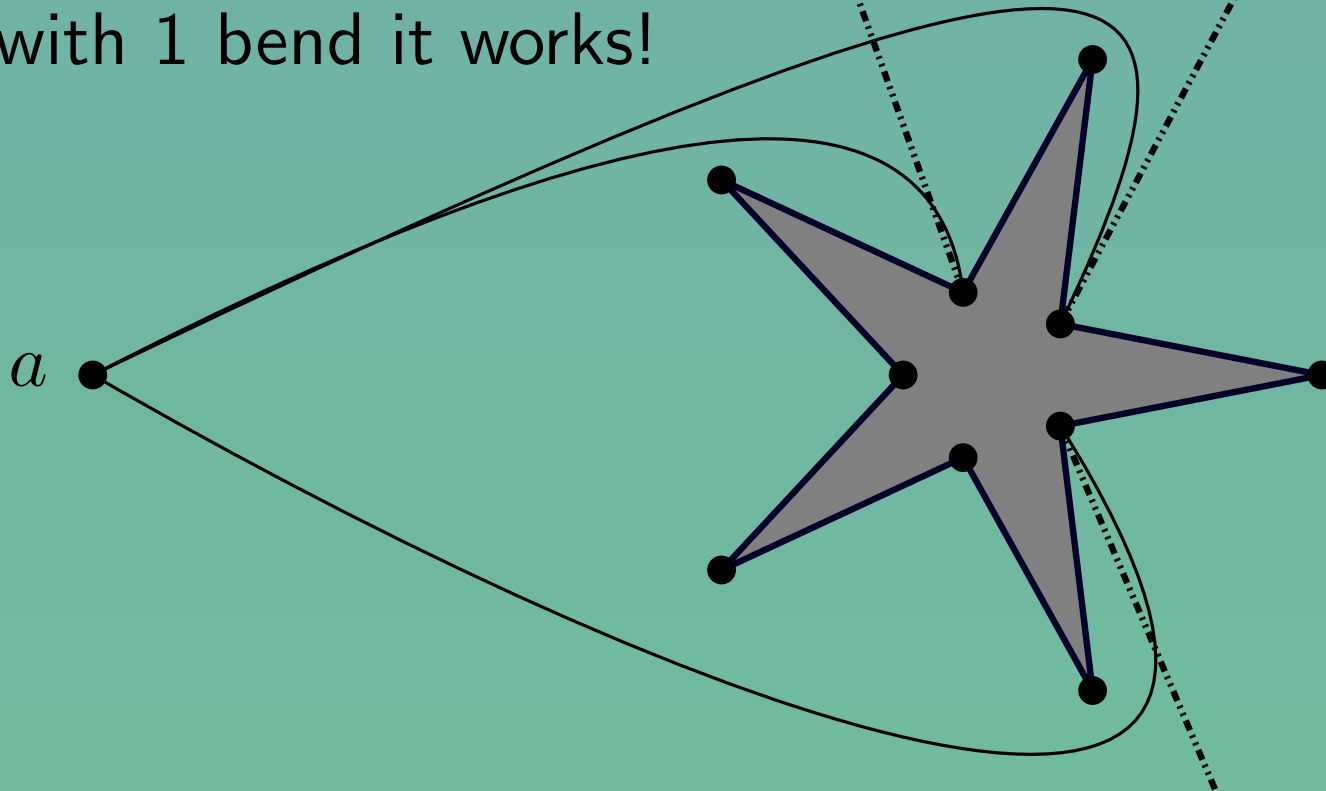


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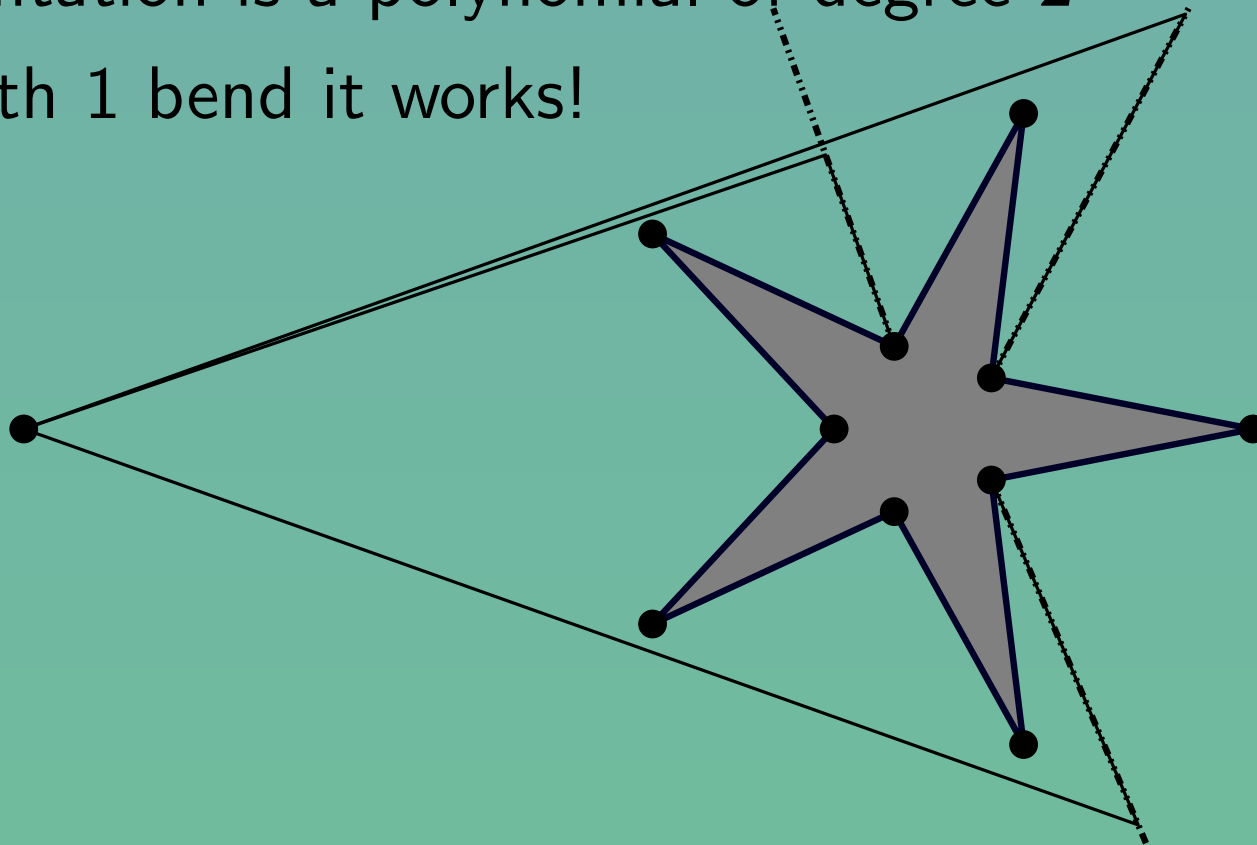


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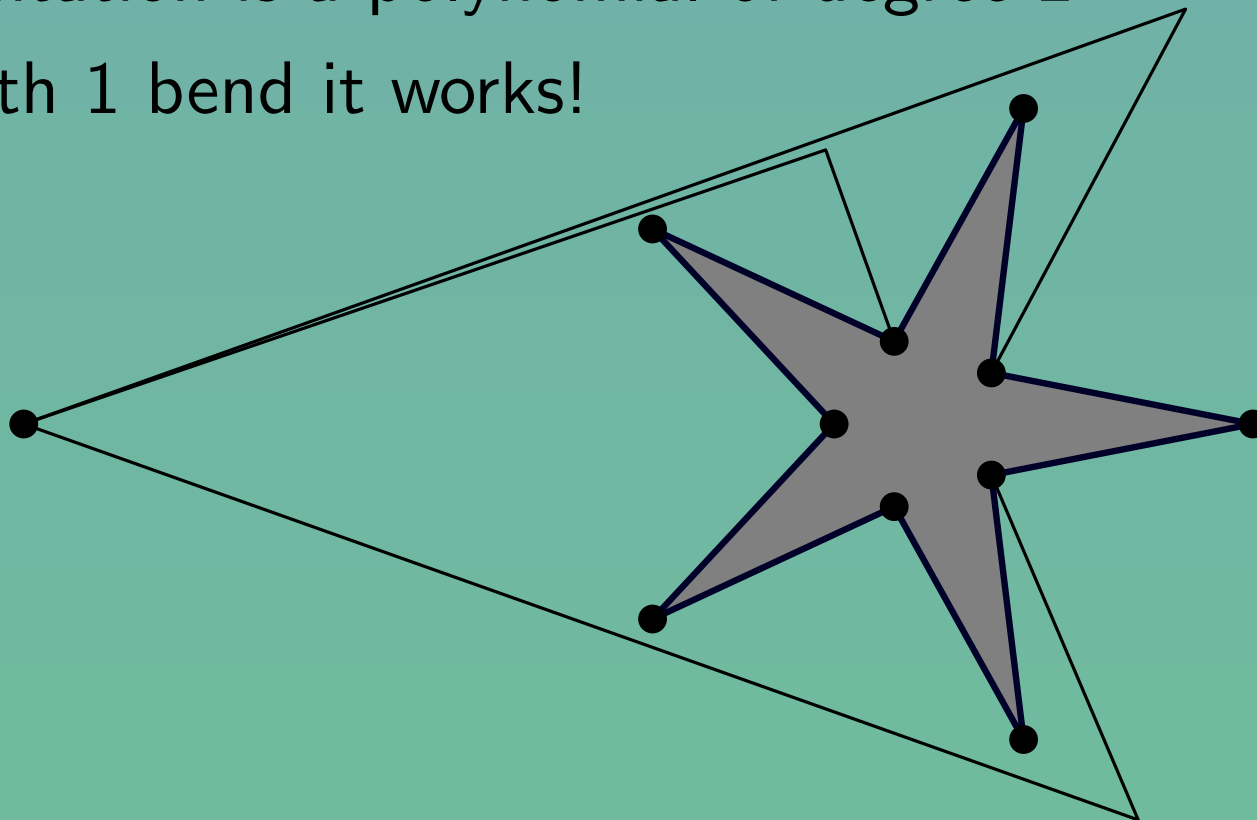


# Extending Stars

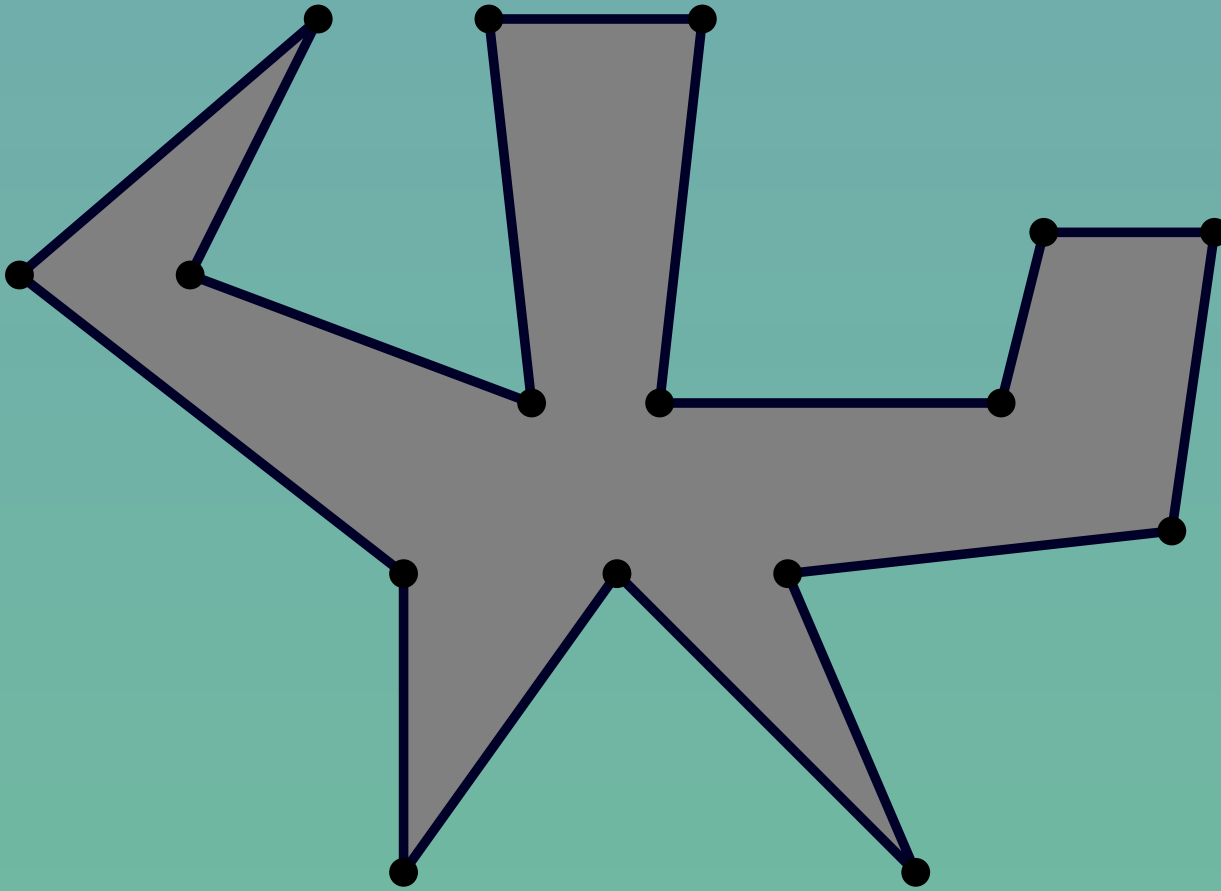
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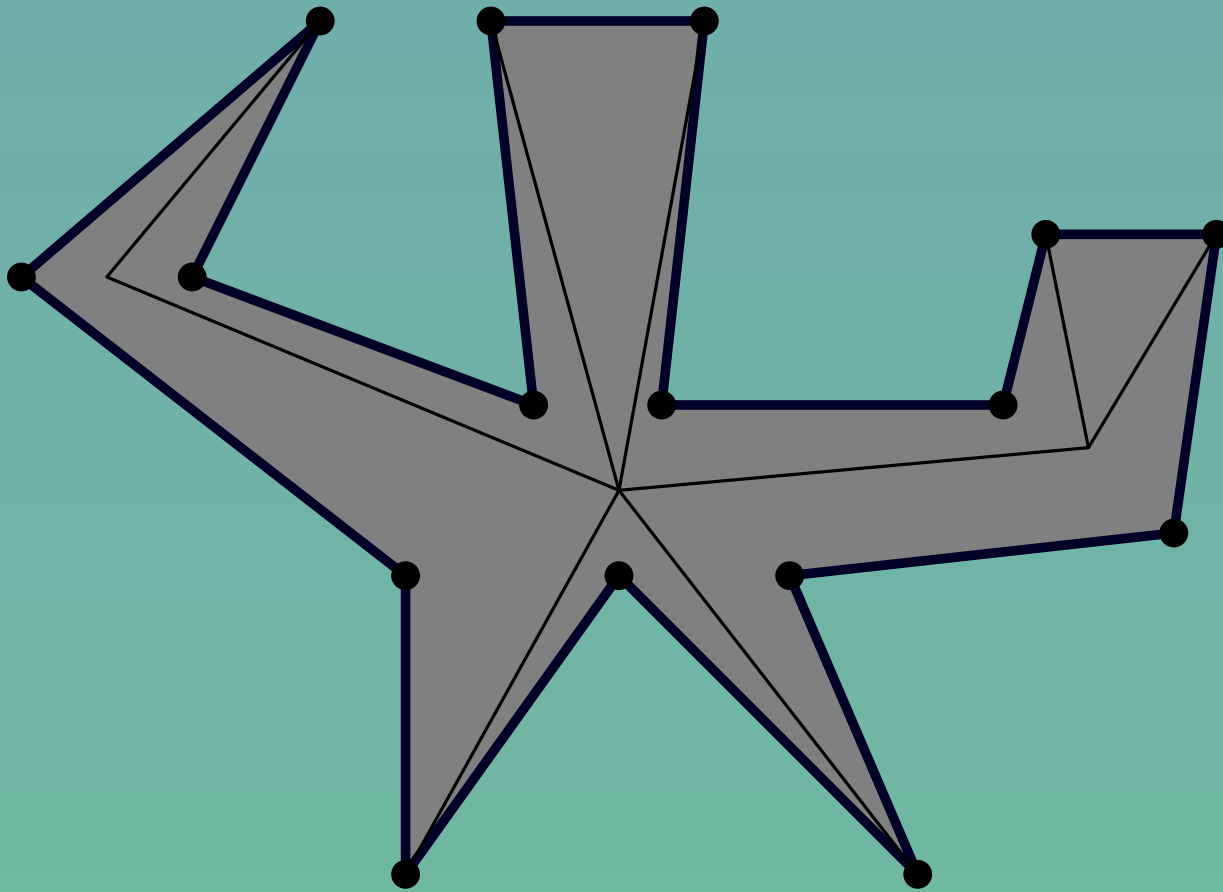
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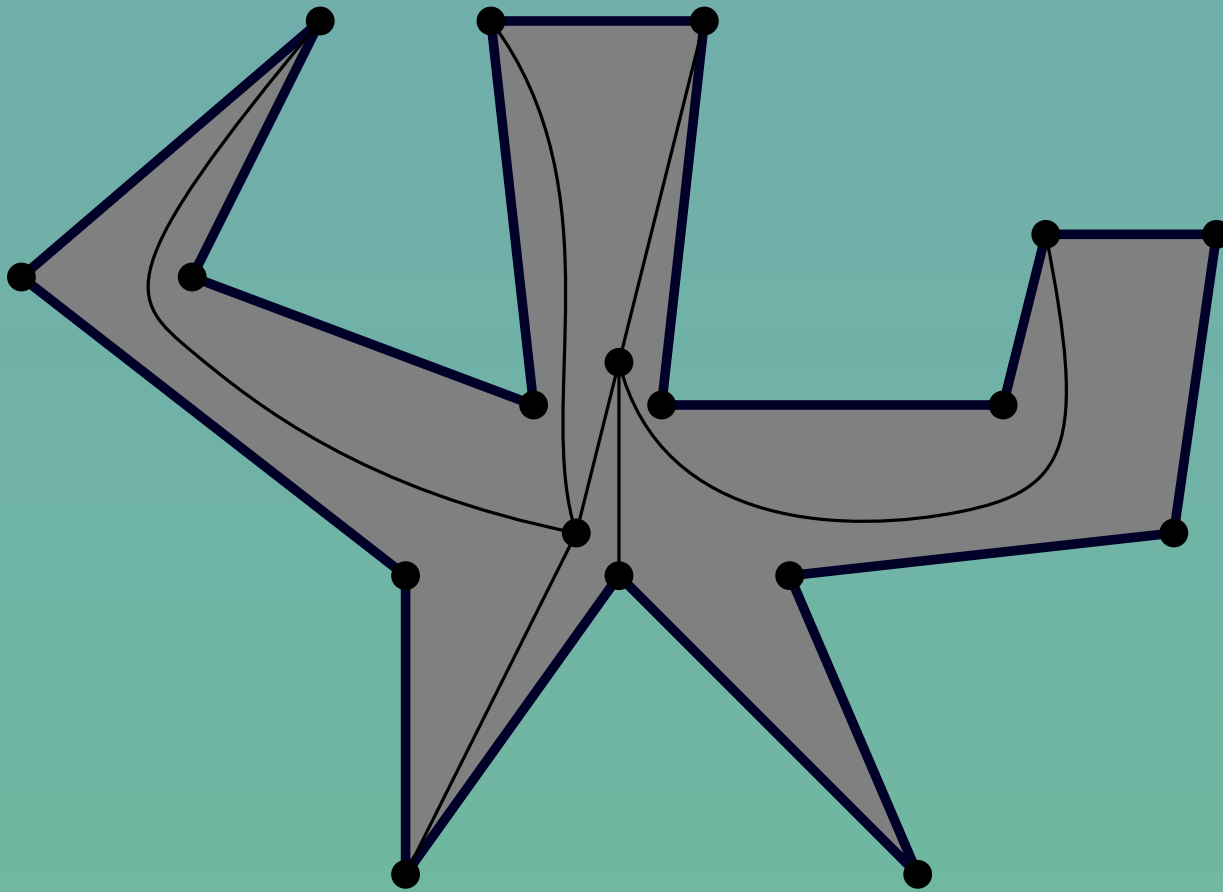
$\beta$ -stars



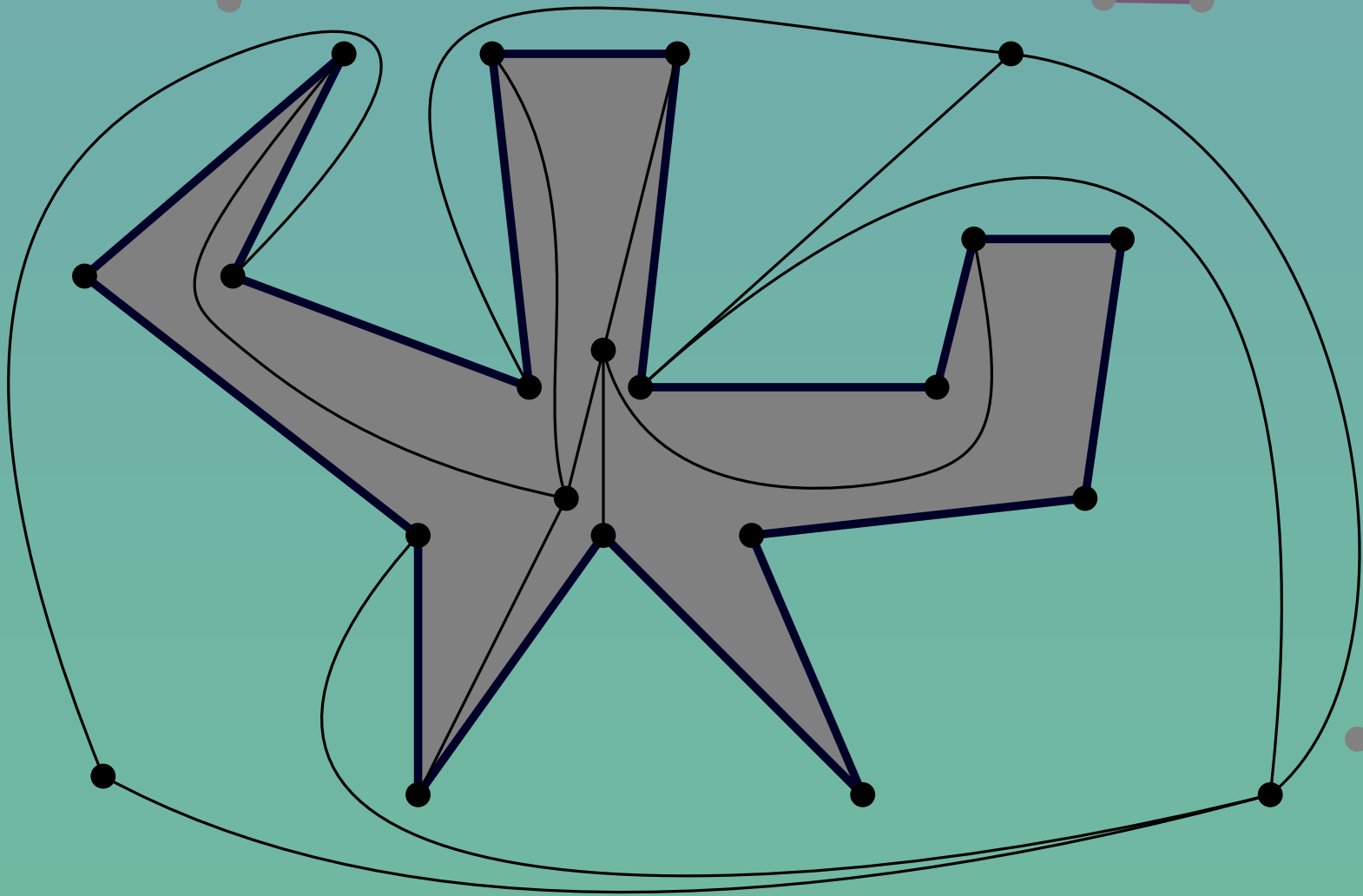
$\beta$ -stars



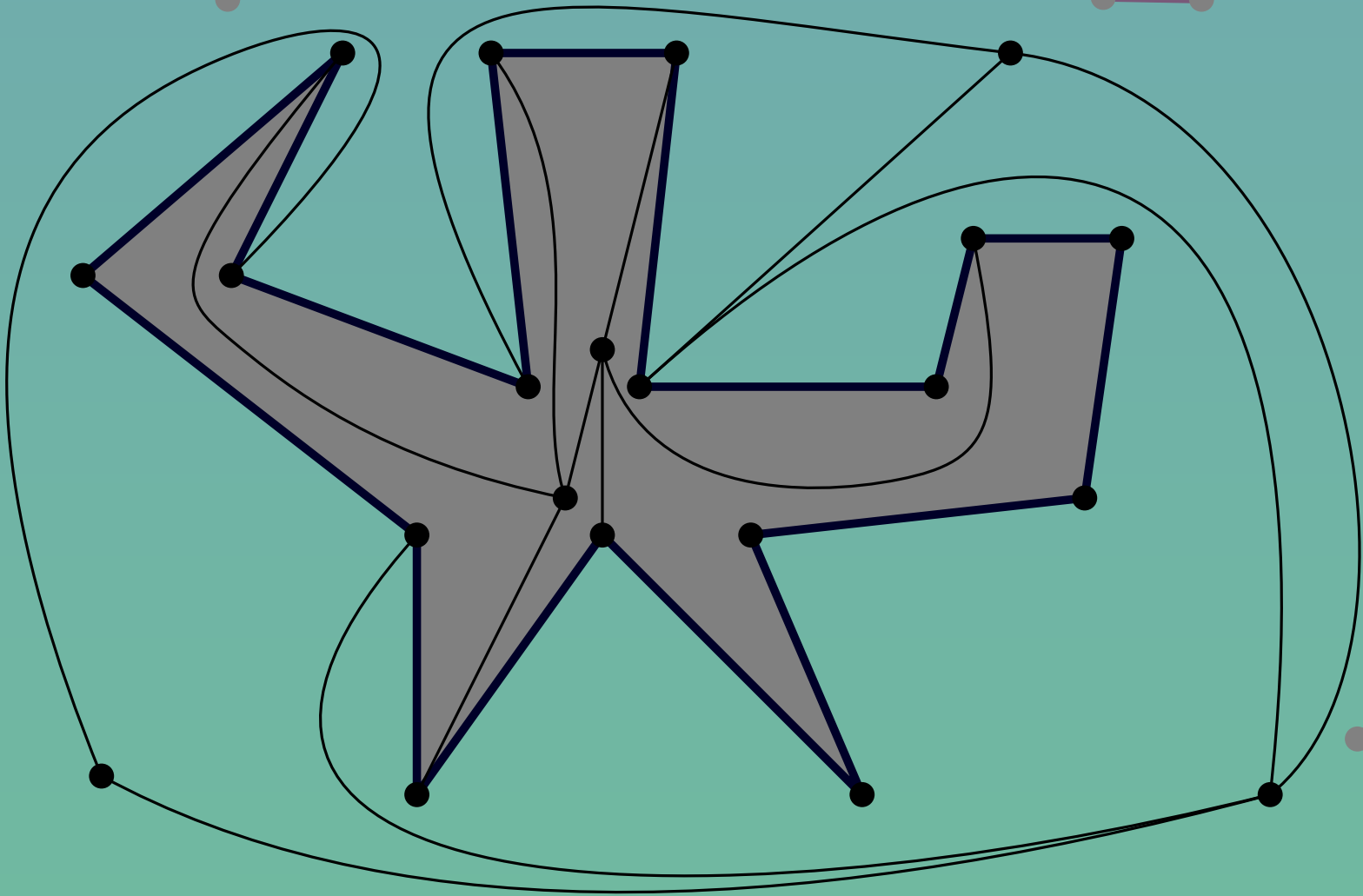
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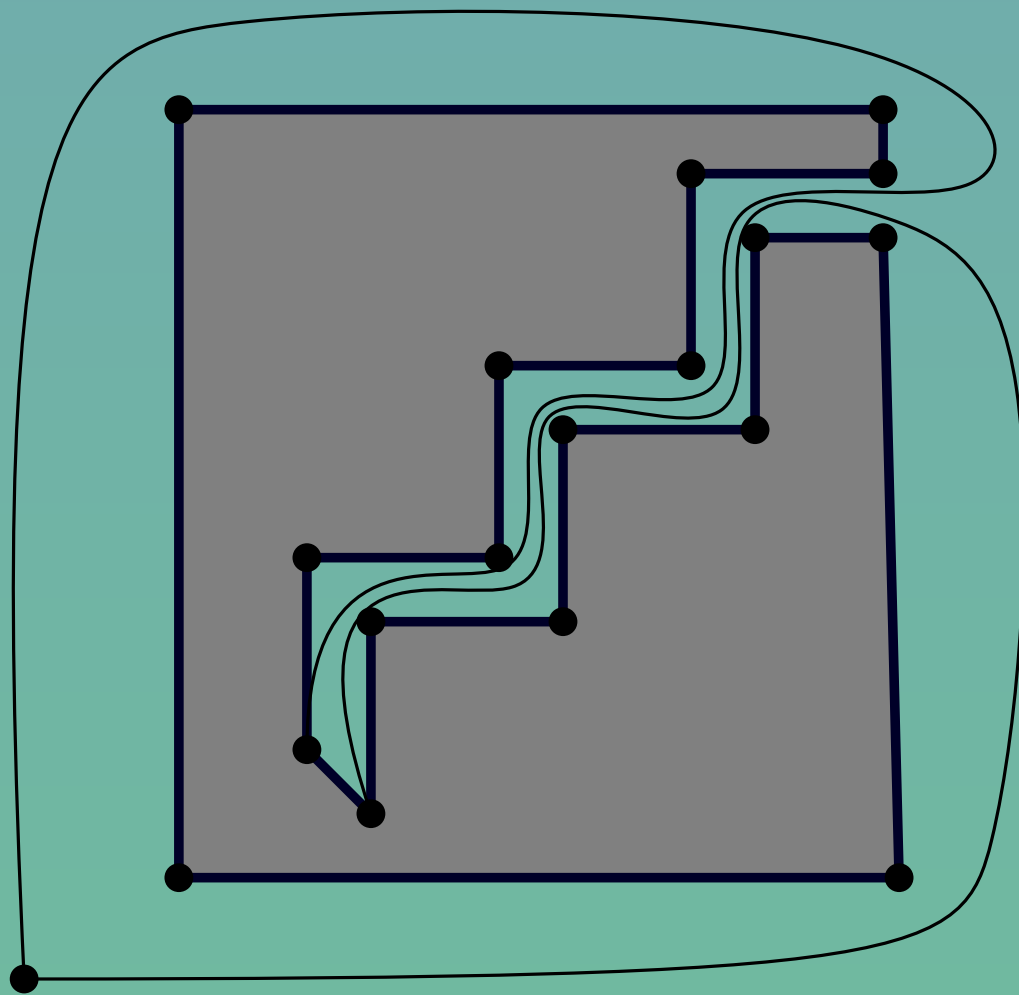


$\beta$ -stars



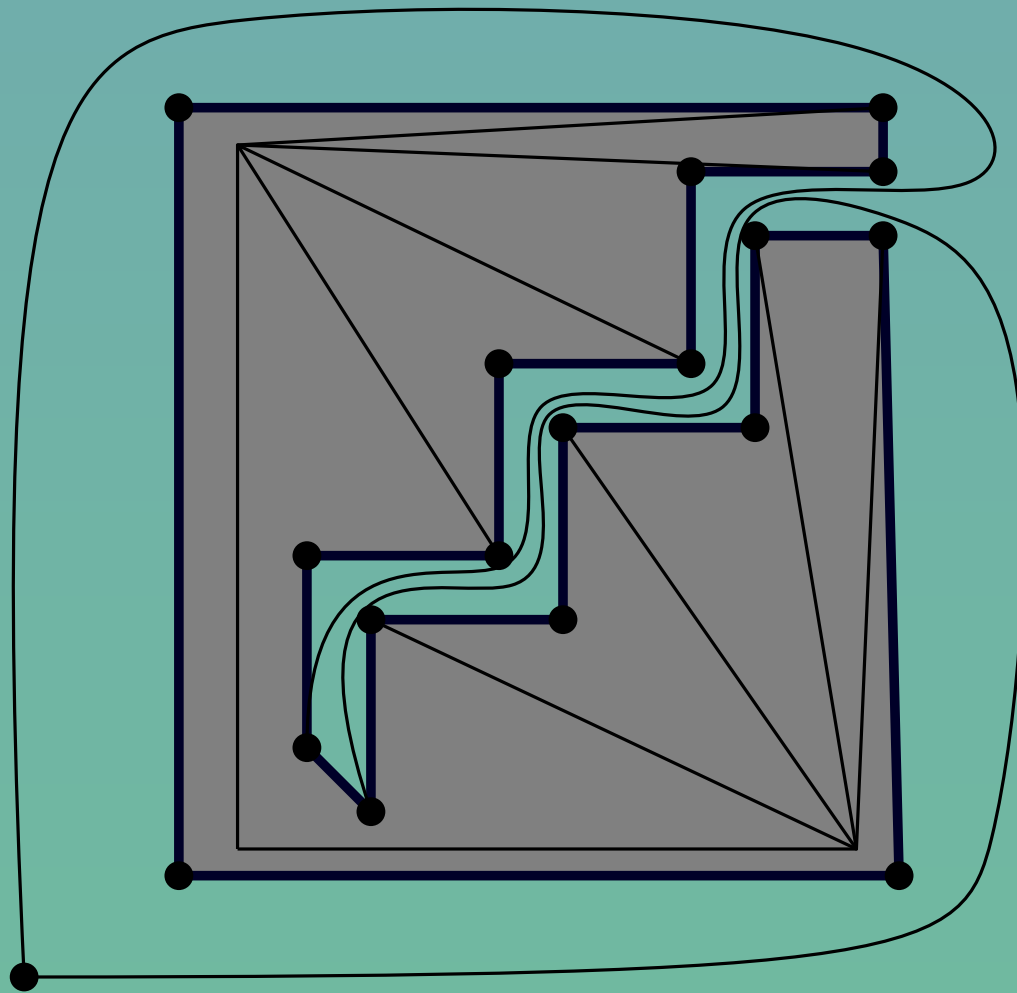
Can each drawing of an outer (inner) face as a  $\beta$ -star be extended with at most  $\beta + 1$  bends per edge?

$\beta$ -stars

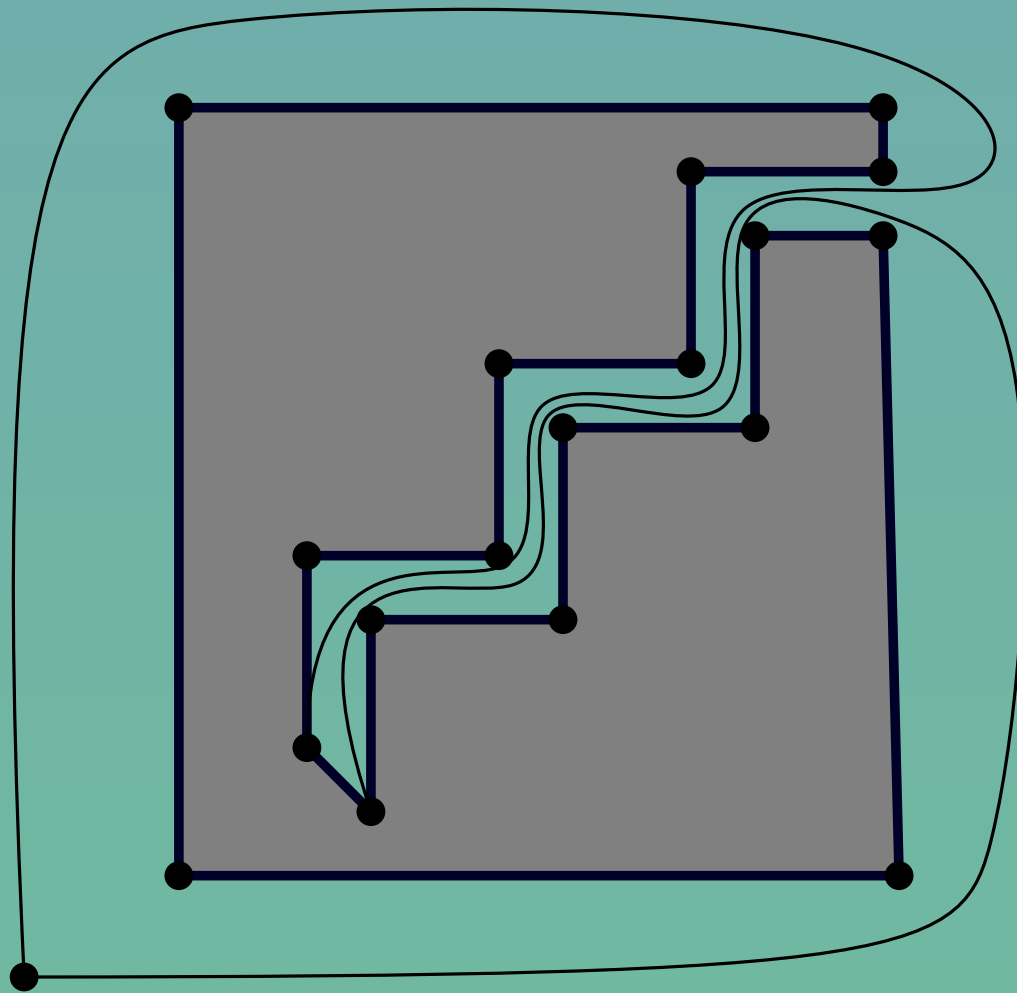




$\beta$ -stars

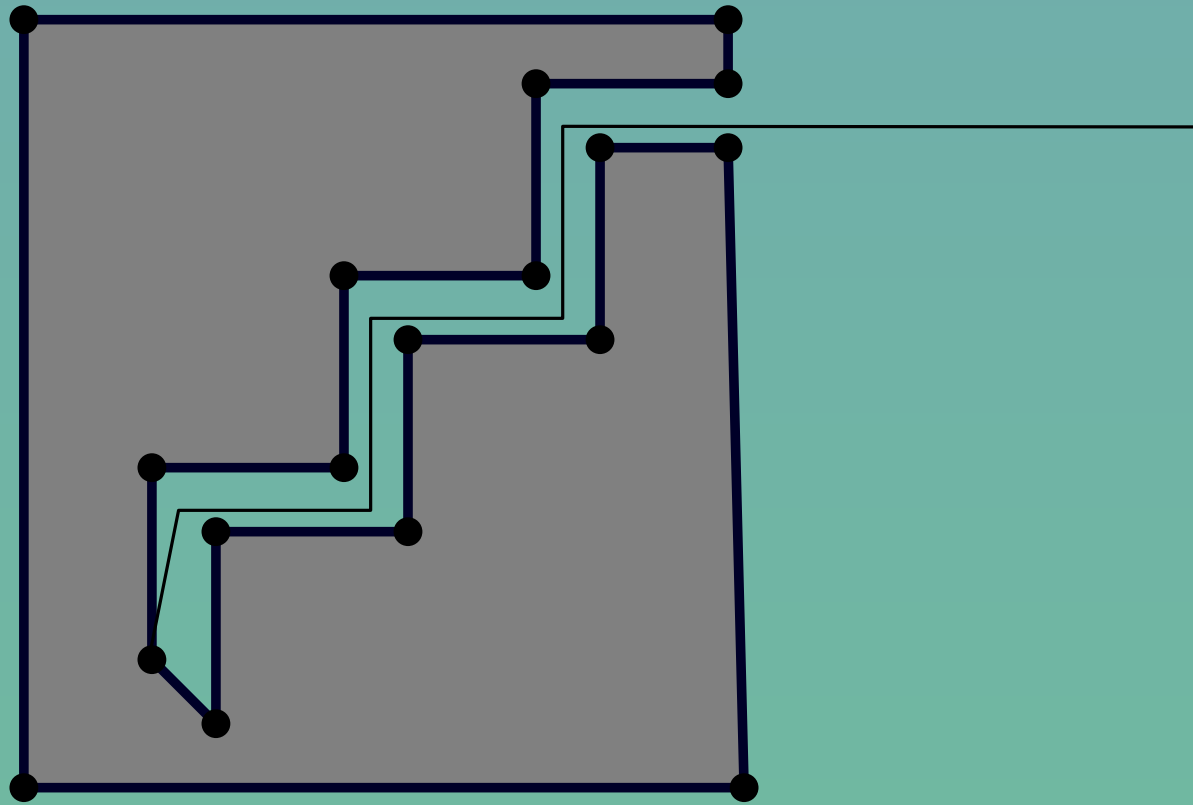


$\beta$ -stars



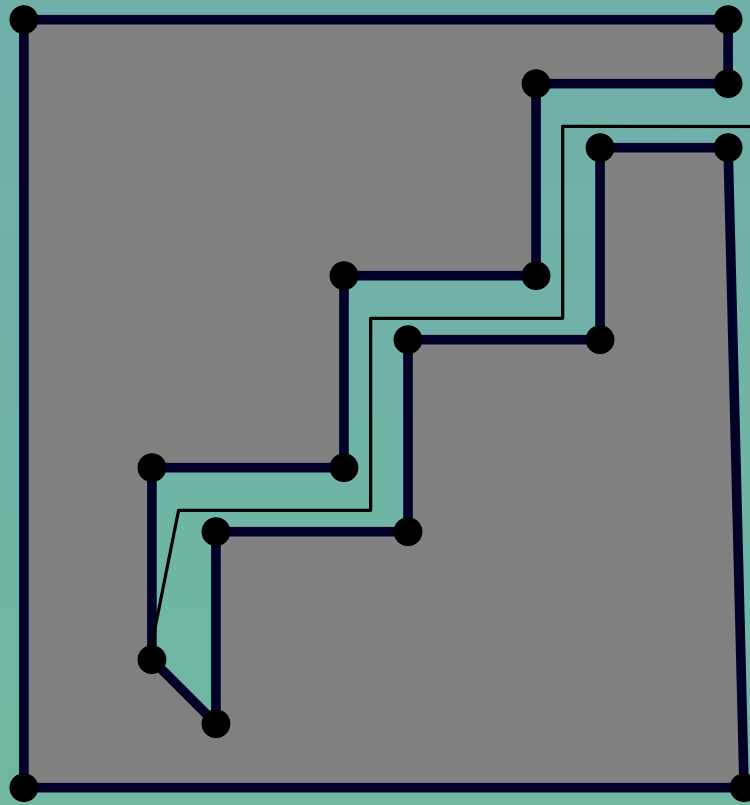
$\beta$ -stars

5-outer-star



$\beta$ -stars

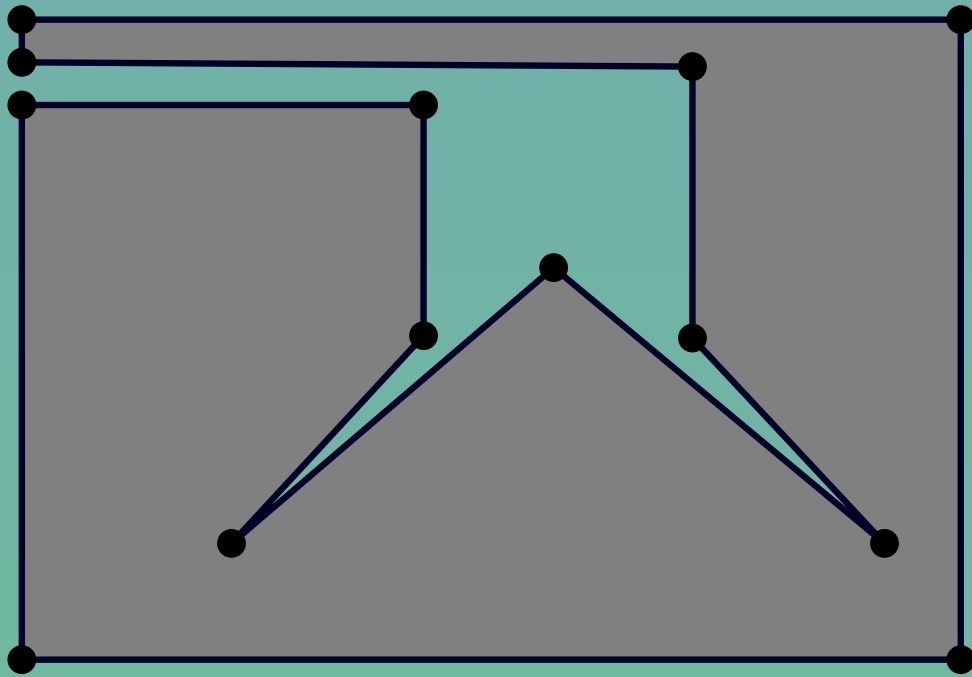
5-outer-star



Can each drawing of an outer (inner) face as a  $\beta$ -(outer)-star be extended with at most  $\beta + 1$  bends per edge?

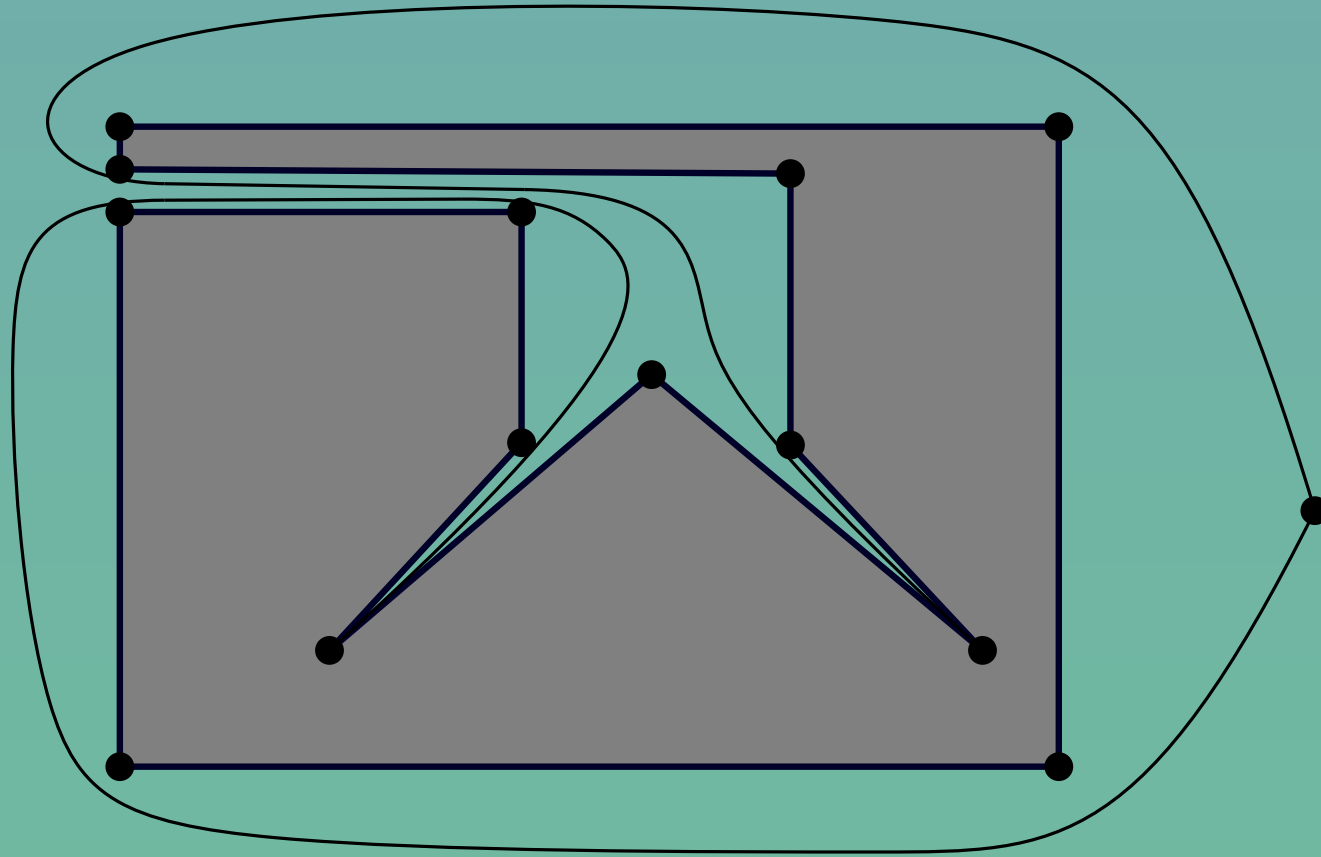
$\beta$ -stars

1-outer-star



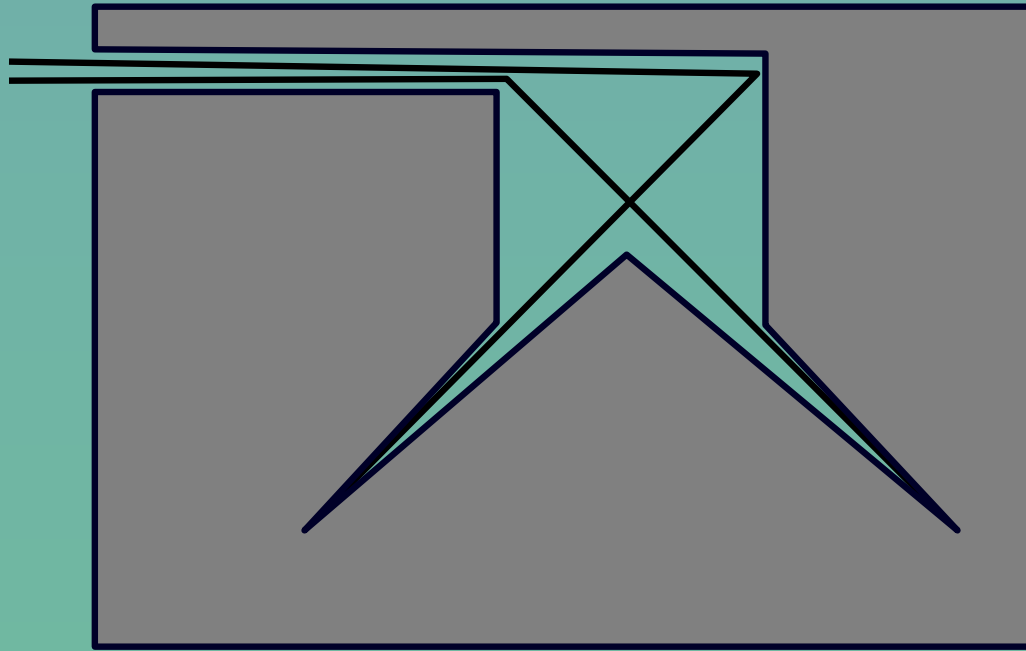
# $\beta$ -stars

1-outer-star, but 3 bends are needed.



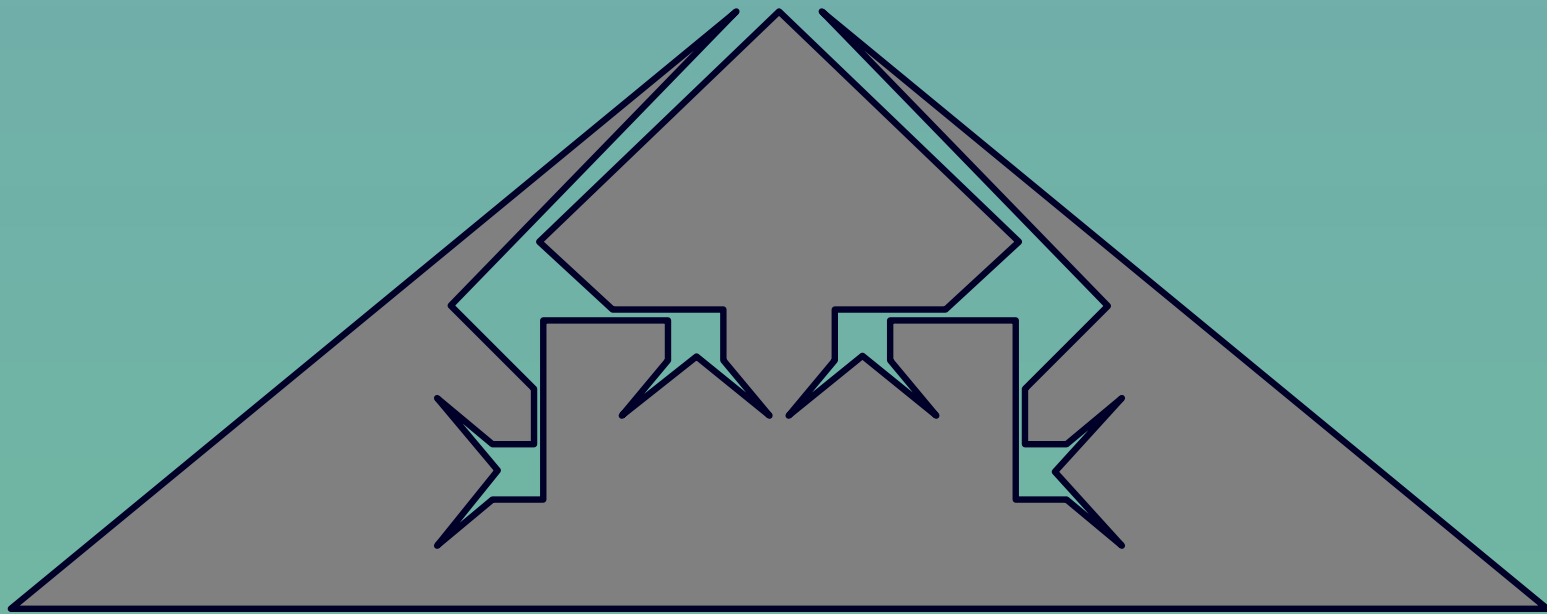
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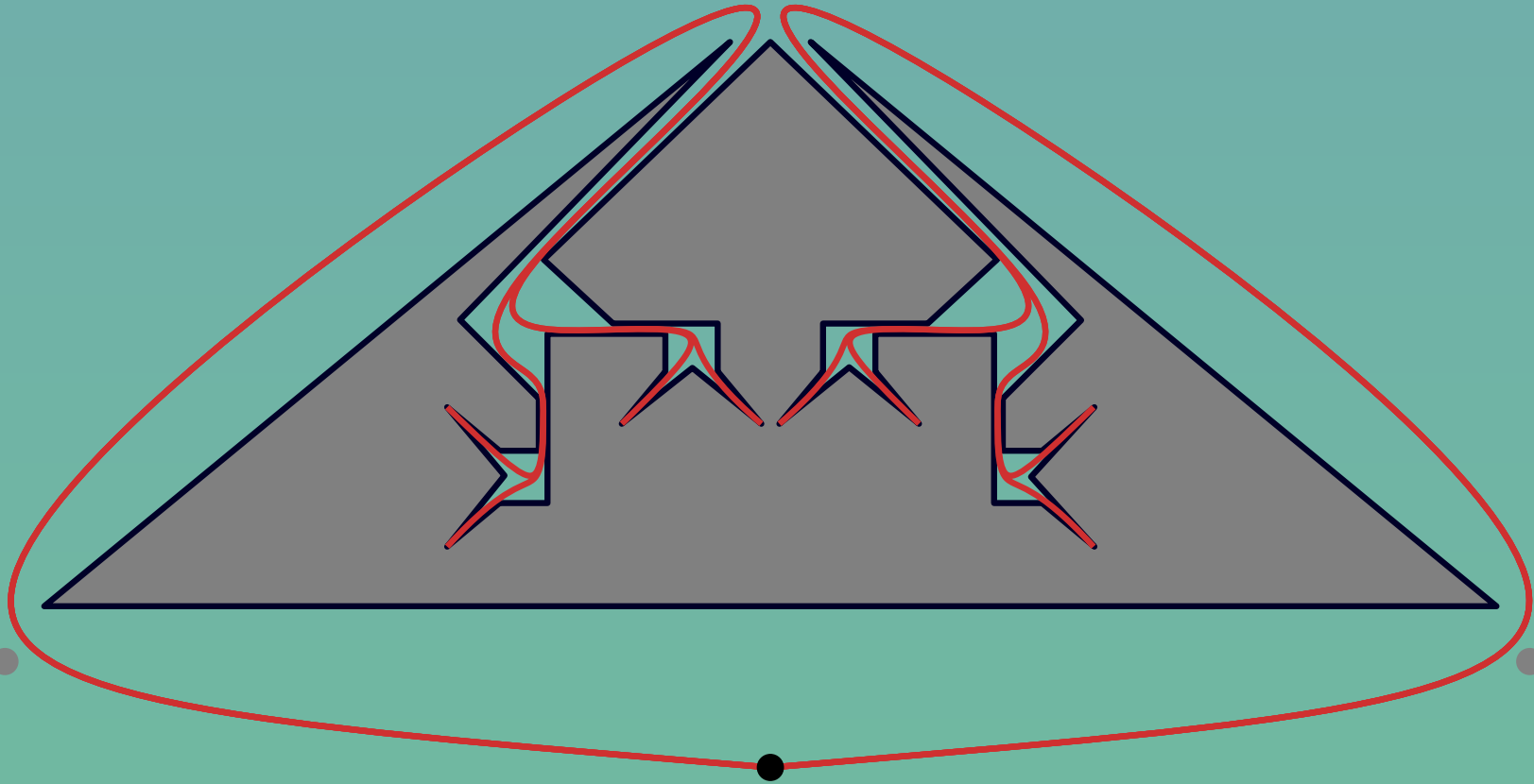
2-outer-star, but 6 bends are needed.





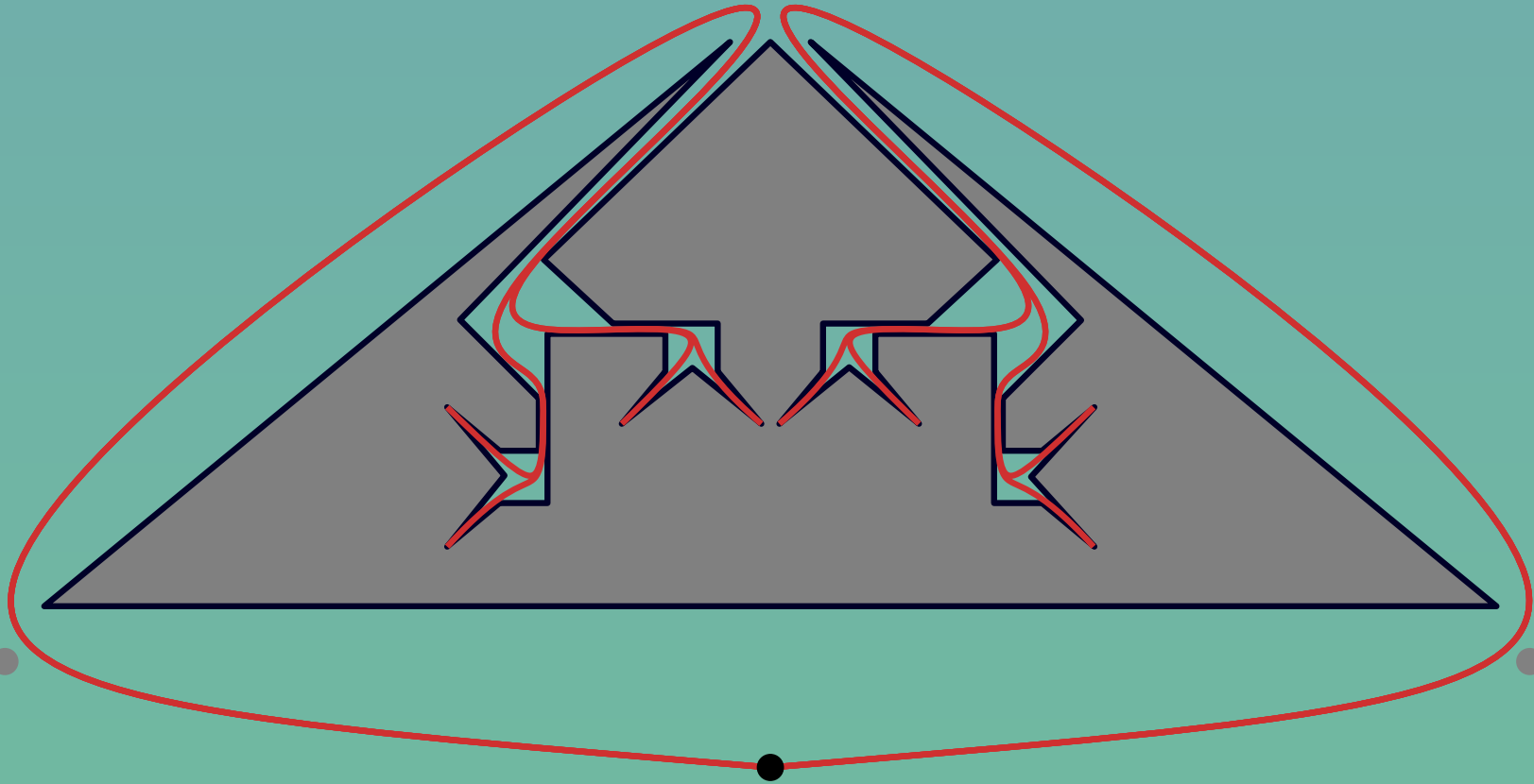
# $\beta$ -stars

2-outer-star, but 6 bends are needed.



$\beta$ -stars

2-outer-star, but 6 bends are needed.



How to get rid of the intersections while adding a minimal number of bends?

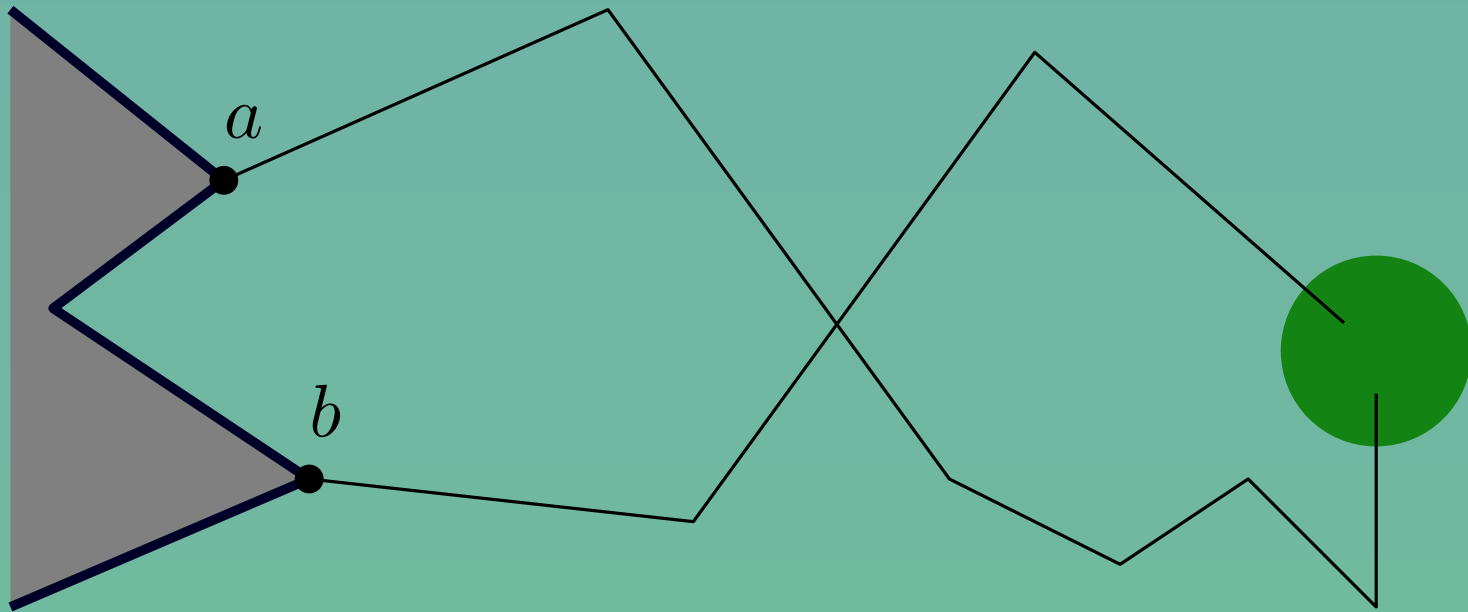
## $\beta$ -stars

Create a set of non-intersecting piecewise-linear curves from each vertex of the face to the kernel / infinity.

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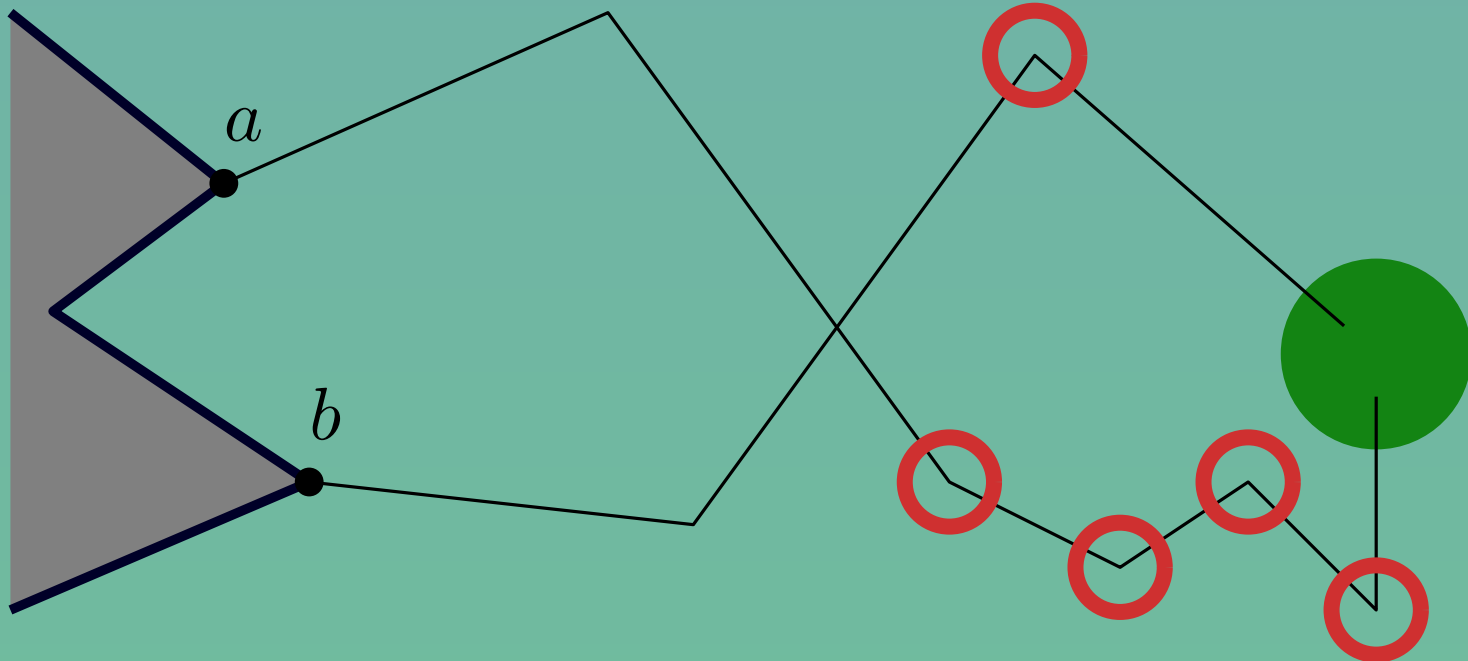
1. remove avoidable intersections



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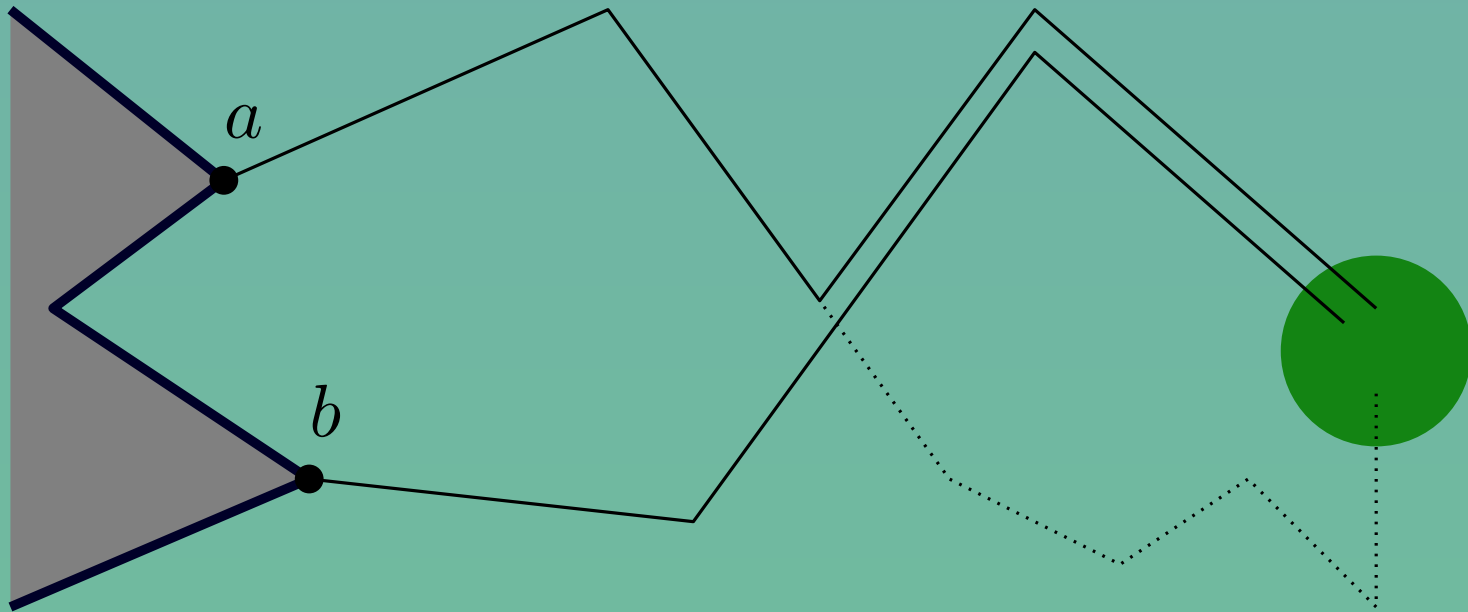
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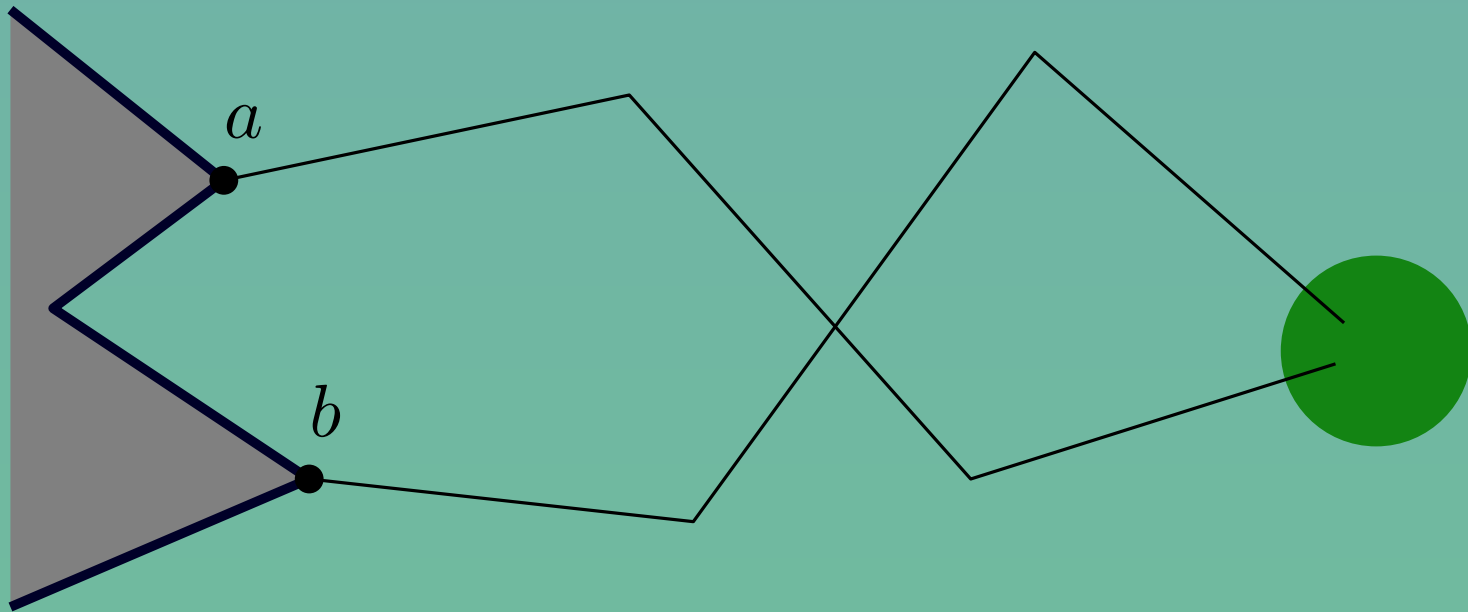
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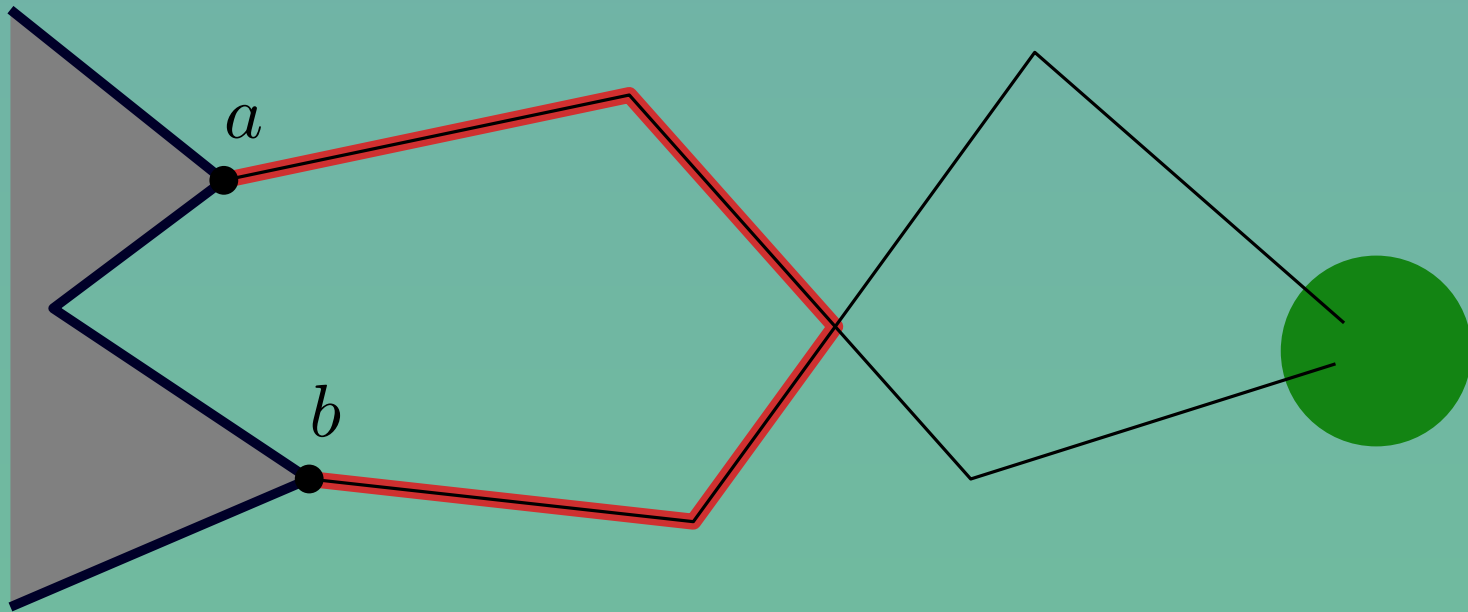
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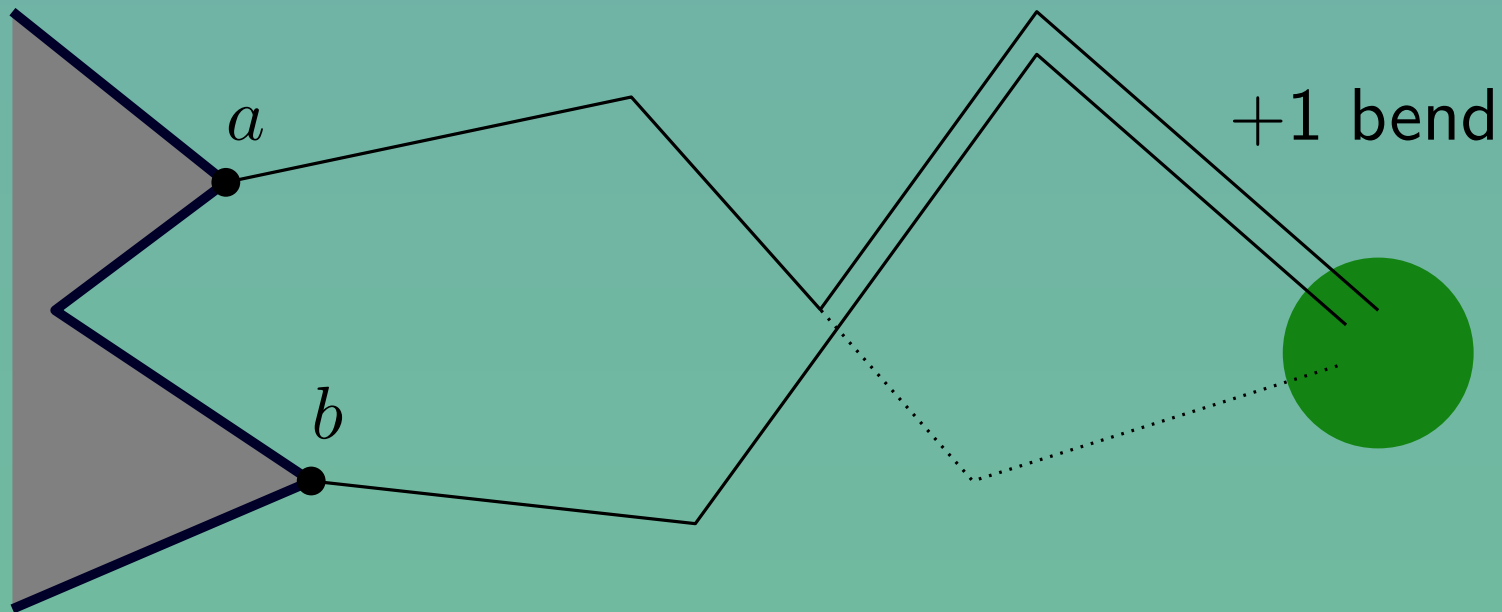




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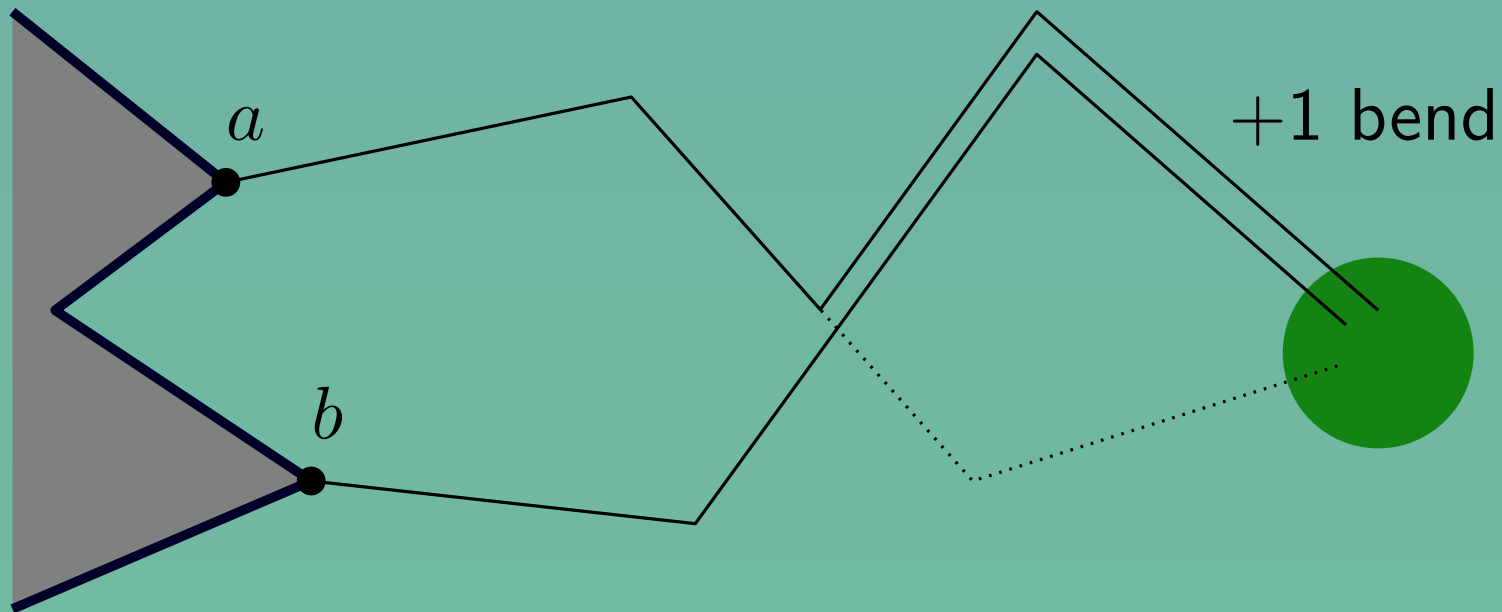
1. remove avoidable intersections
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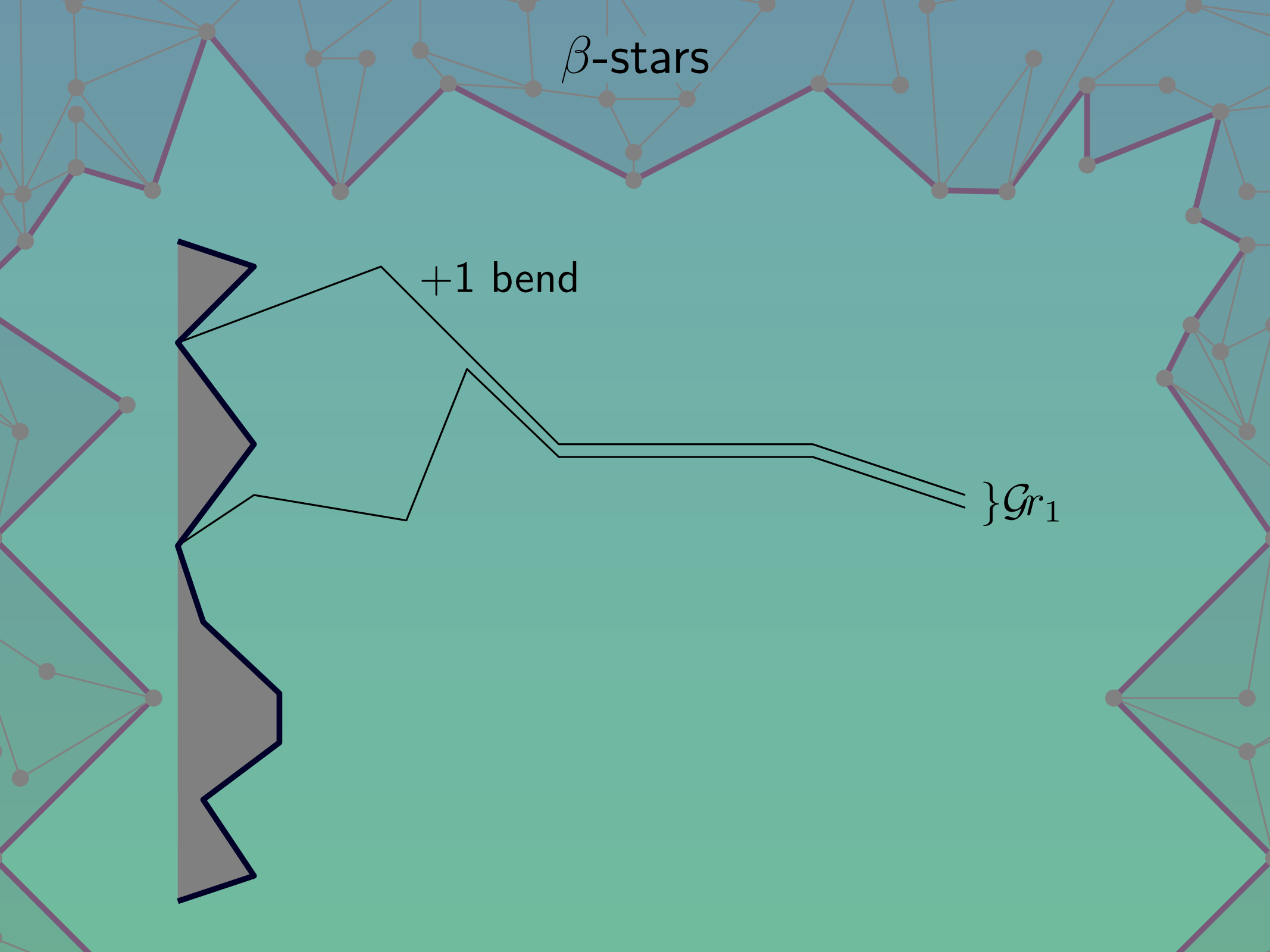
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$\beta$ -stars

+1 bend

$\} Gr_1$



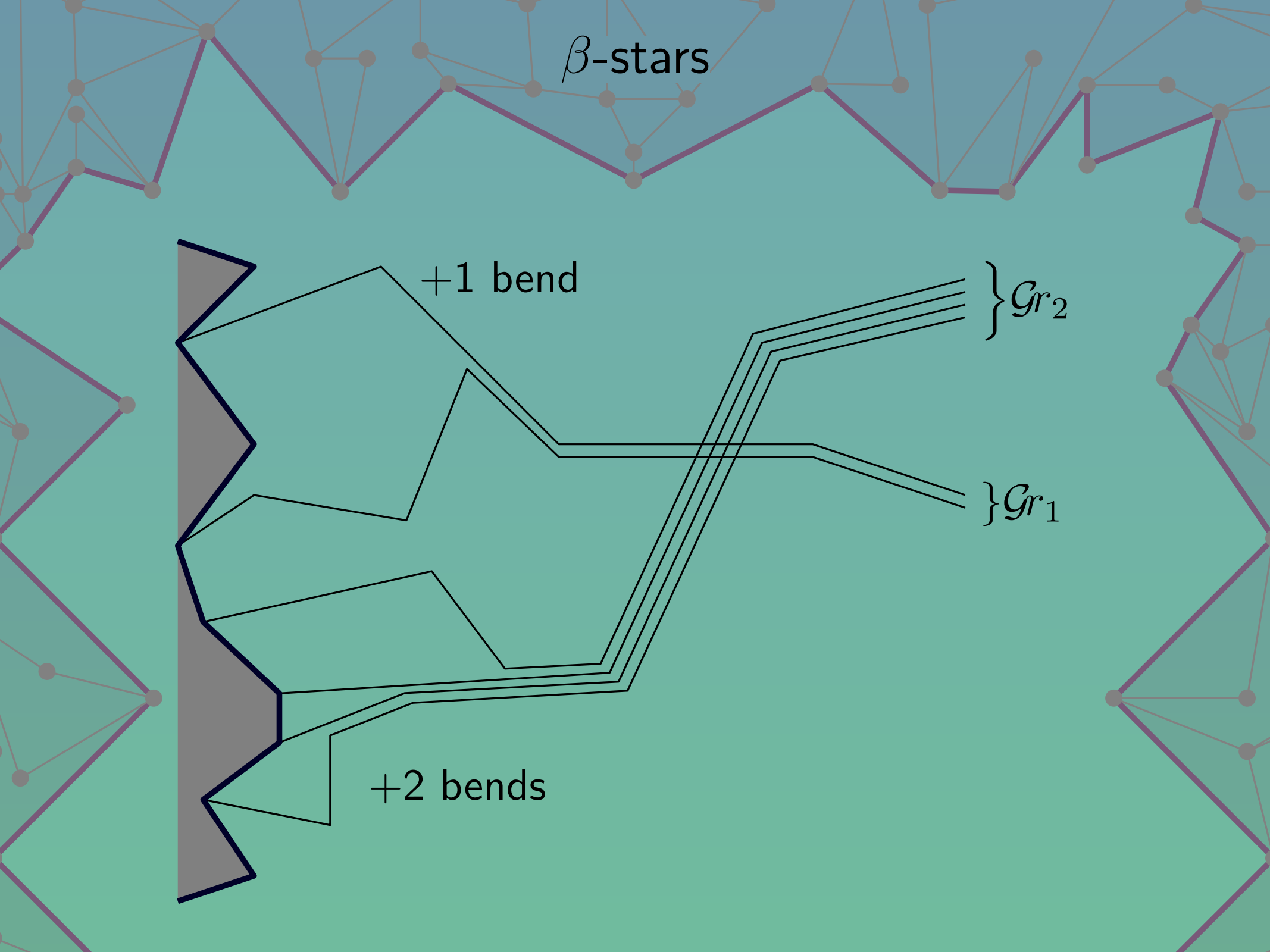
$\beta$ -stars

+1 bend

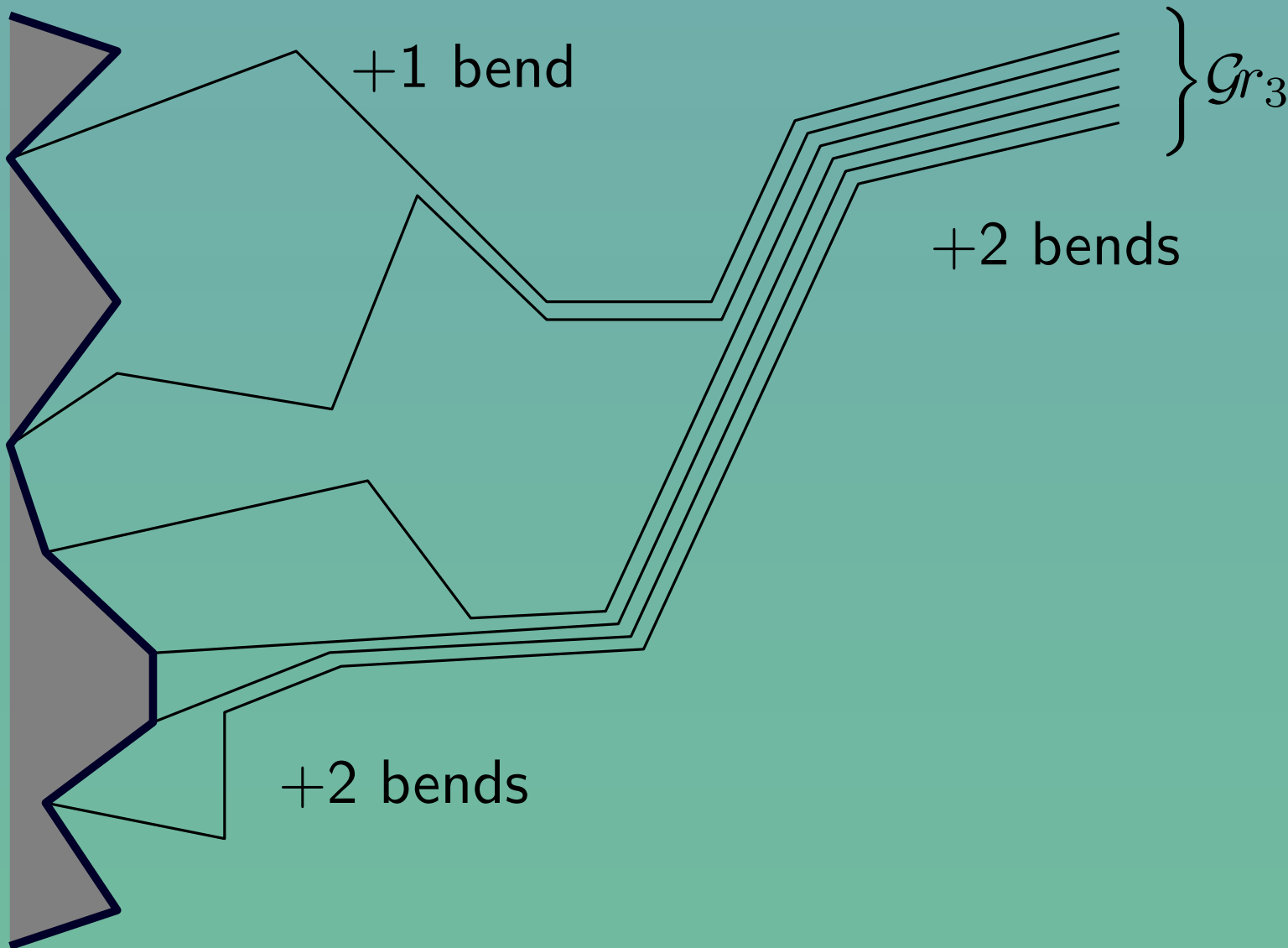
}  $Gr_2$

}  $Gr_1$

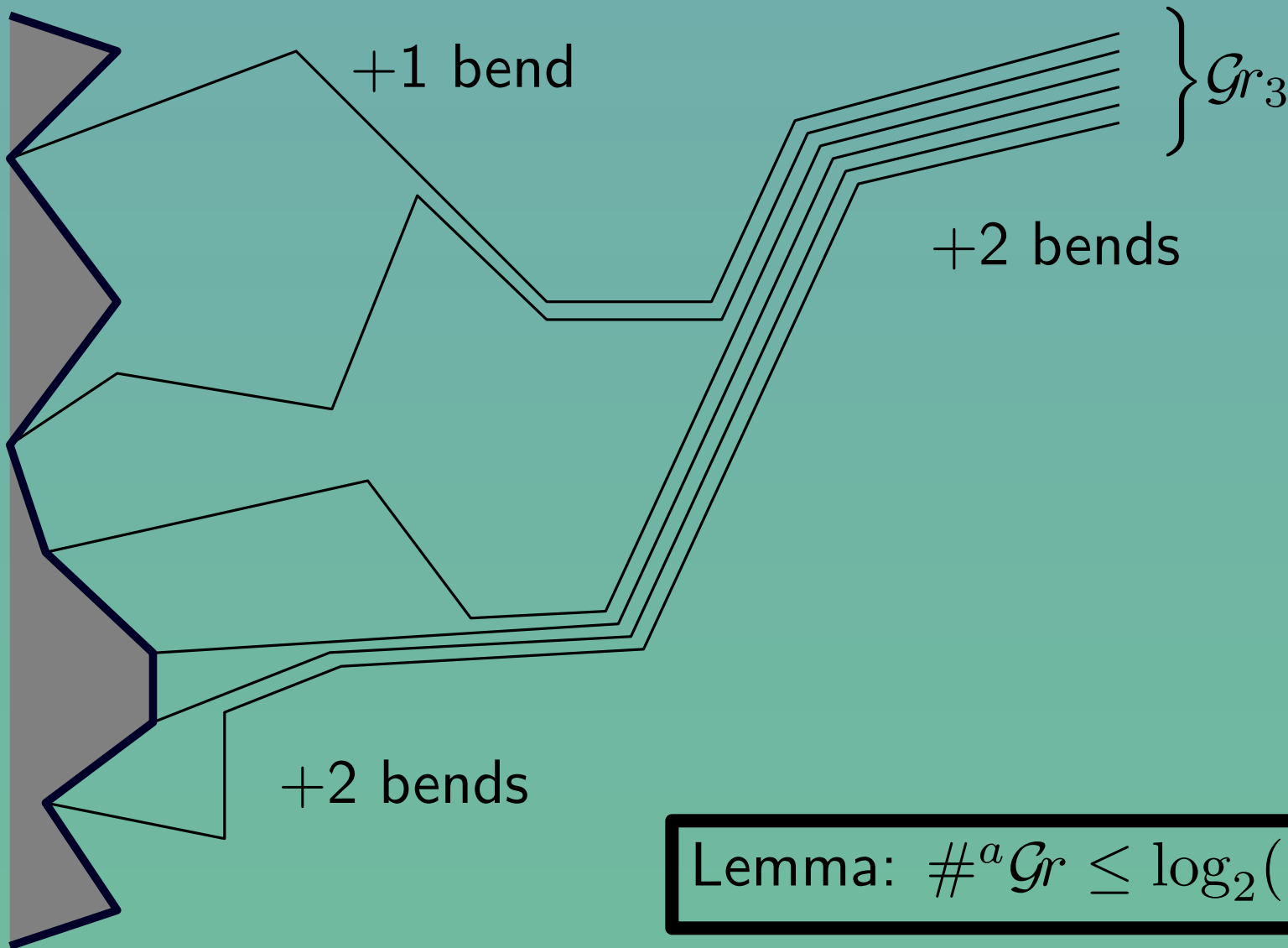
+2 bends



$\beta$ -stars



$\beta$ -stars

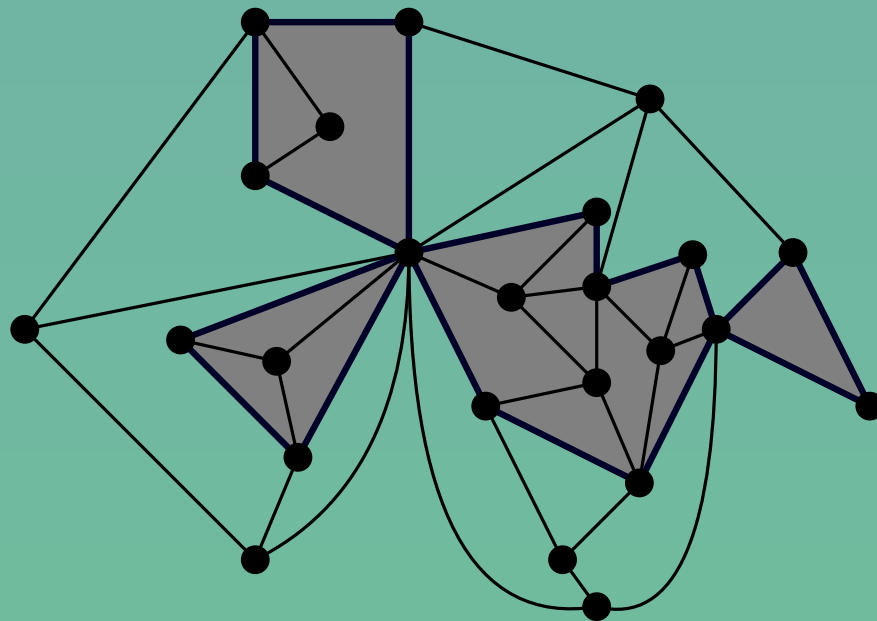


**Lemma:  $\#^a Gr \leq \log_2(|Gr|)$**

$\beta$ -stars

### Theorem

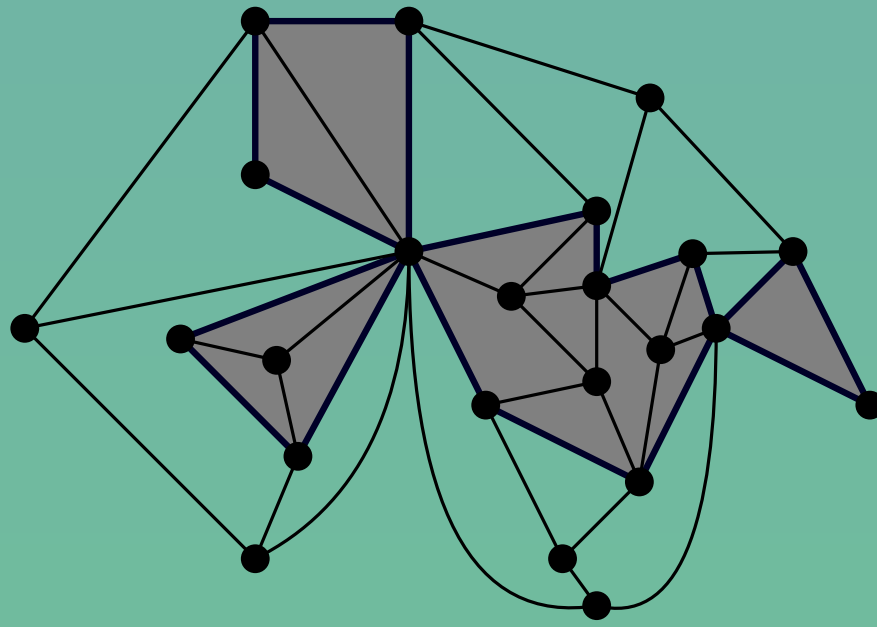
Each instance  $(G, \Gamma_H)$  where  $H$  is an induced connected subgraph of  $G$  allows a  $\min\{|H|/2, \beta + \log_2(|H|) + 1\}$ -bend-extension, where  $\beta$  is the maximum (outer) star complexity of a face in  $\Gamma_H$ . This bound is tight up to an additive constant.



$\beta$ -stars

### Theorem

Each instance  $(G, \Gamma_H)$  where  $H$  is a connected subgraph of  $G$ , allows a  $\min\{|H| + 1, 2\beta + 2 \log_2(|H|) + 3\}$ -bend-extension, where  $\beta$  is the maximum star complexity of a face in  $\Gamma_H$ . This bound is tight up to an additive constant.





## Conclusion

Given a plane graph  $G = (V, E)$ , a drawing of a subgraph  $H$ , extend the drawing using as few bends per edge as possible.

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Given a plane graph  $G = (V, E)$ , a drawing of a subgraph  $H$ , extend the drawing using as few bends per edge as possible.

	outer face	inner face
convex	$0^{[T63],[CEGL12]}$	0-1 & easy test <sup>[MNR16]</sup>
star-shaped	$0^{[HN08]}$	0-1 & ind. of complexity

# Conclusion

Given a plane graph  $G = (V, E)$ , a drawing of a subgraph  $H$ , extend the drawing using as few bends per edge as possible.

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# of bends per edge

$H = (V, \emptyset)$

$3|V| + 2$  [PW01]

-

$72|H|$  [CFGLMS15]

$H$  induced and connected

$\min\{|H|/2, \beta + \log_2(|H|) + 1\}$