Upward Planar Morphs











G. Da Lozzo G. Di Battista F. Frati M. Patrignani V. Roselli





What are we talking about? Morphs

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Transform one drawing into another by moving vertices by preserving at *any time* the properties of the drawings.

Preliminaries

Tools 0000

What are we talking about? Morphs

Transform one drawing into another by moving vertices by preserving at *any time* the properties of the drawings.

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If vertices move at *uniform speed* along straight-line trajectories, the morph is *linear*

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What are we talking about? Morphs

Sometimes linear morphs do not preserve the properties of the input drawings: some intermediate *steps* are necessary

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What are we talking about? Morphs

Sometimes linear morphs do not preserve the properties of the input drawings: some intermediate *steps* are necessary

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Morphs

Linear Morphs

A *linear morph* is a morph such that vertices move at uniform speed along straight-line trajectories.

Morphs

Linear Morphs

A *linear morph* is a morph such that vertices move at uniform speed along straight-line trajectories.

Morphs

A *morph* is a finite sequence of linear morphs, called steps.

The *complexity* of a morphing algorithm is given by the number of intermediate morphing steps.

Morphs Some literature

1914-17: Tietze, Smith, Veblen, [...]

existential proofs for polygons

1944-83: Cairns, Thomassen, [...]

o existential proofs for triangulations and planar graphs

2006: Lubiw et al., SODA

o morphs of orthogonal drawings in a polynomial number of steps

2013: Alamdari et al., SODA

 first algorithm for morphing general planar graph drawings in polynomially many steps

2014: Angelini et al., ICALP

o optimal (linear) algorithm for planar graphs

2014: Barrera-Cruz et al., GD

polynomial algorithm for Schnyder drawings of triangulations

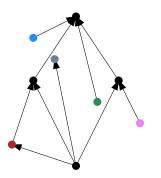
2015: Angelini et al., SoCG

optimal (linear) algorithm for convex drawings

2018: van Goethem & Verbeek, SoCG

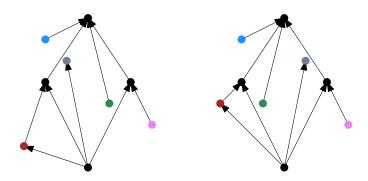
polynomial algorithm for orthogonal drawings

What are we talking about? Upward Plane Graphs



- natural extension of planarity to directed graphs
- edges drawn upward
- vertices with no incoming edges are sources
- vertices with no outgoing edges are sinks

What are we talking about? Upward Plane Graphs



Upward Planar Morphs are defined for *upward-equivalent* drawings.

Upward equivalence is a *necessary condition*.

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Plane st-Graphs

Upward Drawings Some literature

1988: Di Battista & Tamassia

drawing algorithms

1992: Di Battista et al.

o area requirements

1994-98: Bertolazzi et al.

o drawing algorithms for triconnected, upward planarity

test for single-source

2002: Garg & Tamassia

o upward planarity test

...

Upward Planar Morphs

Our results

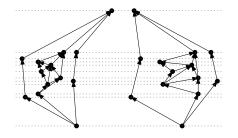
Given two upward-equivalent planar straight-line drawings, there *always exists* a morph between them such that *all the intermediate drawings* of the morph are upward planar and straight-line.

Upward Planar Graphs	$O(n^2)$ – $\Omega(n)$ steps
Reduced Upward Planar Graphs	$\Theta(n)$ steps
Planar st-Graphs	O(n) steps
Reduced Planar st-Graphs	$\Theta(1)$ steps

Definitions & Tools

Left-to-Right equivalence

Two upward-equivalent drawings Γ_0 and Γ_1 of a graph are LR-equivalent if any horizontal line intersects the same sequence of vertices and/or edges in both drawings.



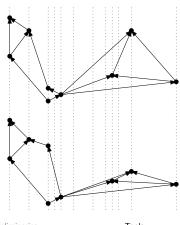
For any two LR-equivalent drawings of G, there exists a 1-step upward planar morph between them.

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Bottom-to-Top equivalence

Two upward-equivalent drawings Γ_0 and Γ_1 of a graph are BT-equivalent if any vertical line intersects the same sequence of vertices and/or edges in both drawings.



For any two LR-equivalent drawings of G, there exists a 1-step upward planar morph between them.

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Two upward-equivalent planar drawings Γ_0 and Γ_1 of G are an HVH-pair if there exist Γ_0' and Γ_1' such that:





Two upward-equivalent planar drawings Γ_0 and Γ_1 of G are an HVH-pair if there exist Γ_0' and Γ_1' such that:

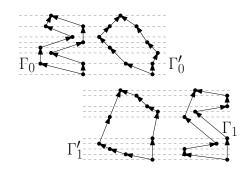
ullet Γ_0 and Γ_0' are LR-equivalent





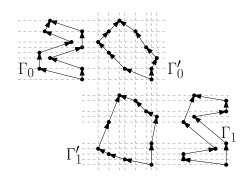
Two upward-equivalent planar drawings Γ_0 and Γ_1 of G are an HVH-pair if there exist Γ_0' and Γ_1' such that:

- ullet Γ_0 and Γ_0' are LR-equivalent
- ullet Γ_1 and Γ_1' are LR-equivalent



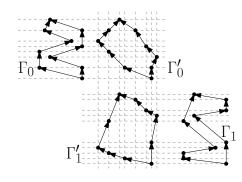
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- Γ_0 and Γ_0' are LR-equivalent
- ullet Γ_1 and Γ_1' are LR-equivalent
- \bullet Γ_0' and Γ_1' are BT-equivalent



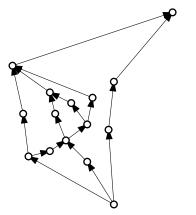
Two upward-equivalent planar drawings Γ_0 and Γ_1 of G are an HVH-pair if there exist Γ_0' and Γ_1' such that:

- ullet Γ_0 and Γ_0' are LR-equivalent
- Γ_1 and Γ_1' are LR-equivalent
- \bullet Γ_0' and Γ_1' are BT-equivalent

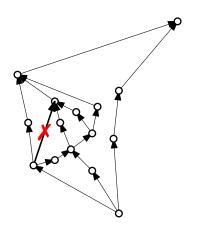


Lemma

For any two upward planar drawings that form an HVH pair, there exists a 3-step upward planar morph between them.



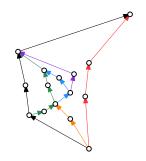
• G is a DAG with exactly one source and one sink, and they lie on the outer face



- *G* is a DAG with exactly one source and one sink, and they lie on the outer face
- reduced = no transitive edges

Theorem

For any two upward-equivalent planar drawings of a reduced st-graph G, there exists a 3-step upward planar morph between them.



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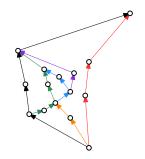
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Theorem

For any two upward-equivalent planar drawings of a reduced st-graph G, there exists a 3-step upward planar morph between them.

Proof: we show that they are an HVH-pair



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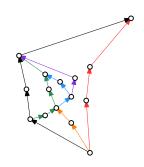
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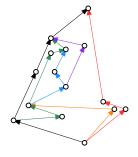
Theorem

For any two upward-equivalent planar drawings of a reduced st-graph G, there exists a 3-step upward planar morph between them.

Proof: we show that they are an HVH-pair



- assume G biconnected (or augment it)
- \bullet exploit an ear-decomposition for constructing Γ_0' and Γ_1'

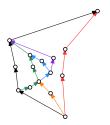


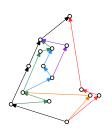
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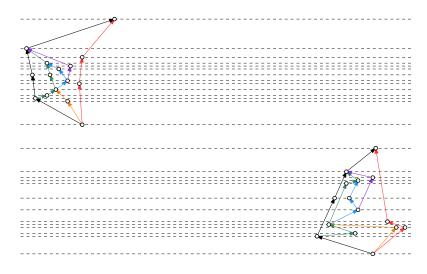
Reduced Plane st-Graphs Construction of Γ'_0 and Γ'_1





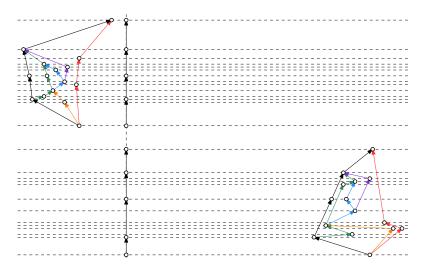
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Construction of Γ_0' and Γ_1'



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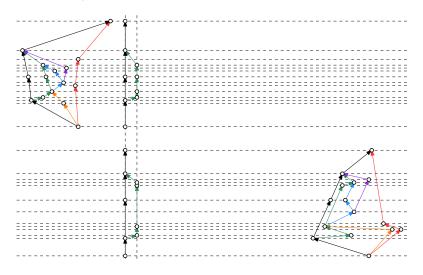
Construction of Γ_0' and Γ_1'



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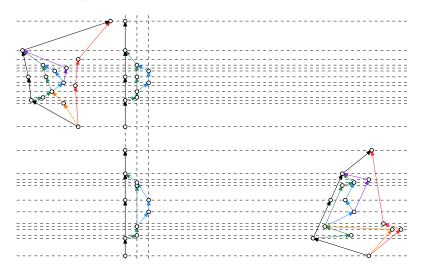
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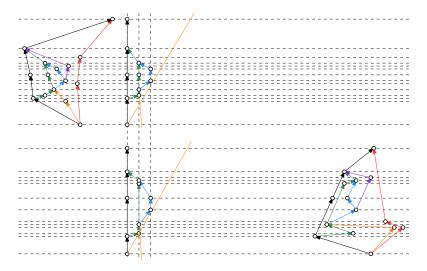
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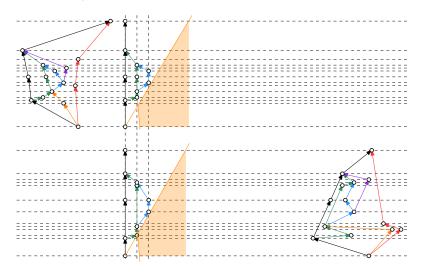
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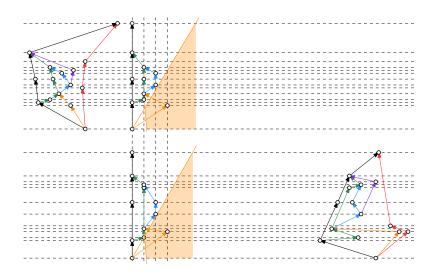
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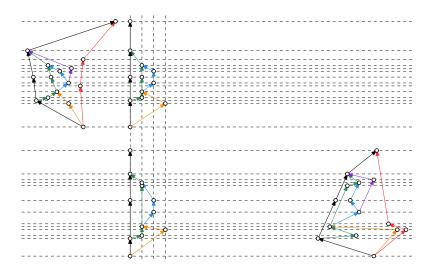
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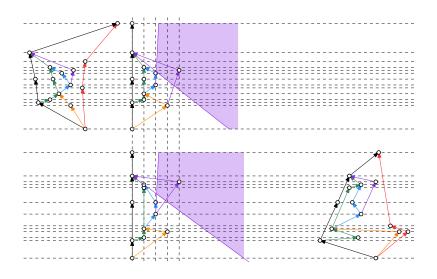
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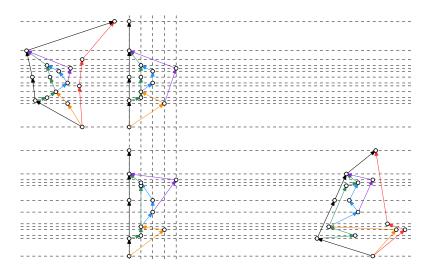
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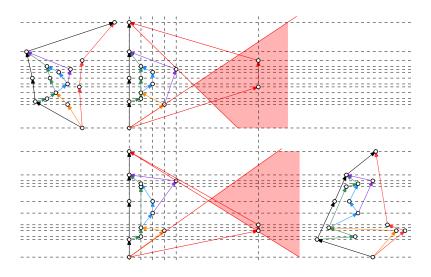
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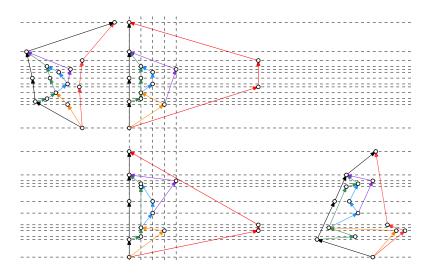
Construction of Γ_0' and Γ_1'



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Construction of Γ_0' and Γ_1'



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Theorem

For any two upward-equivalent planar drawings of an n-vertex plane st-graph G, there exists a O(n)-step upward planar morph between them.

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Proof

ullet augment G to be a maximal st-graph, call it G^+

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Proof

- ullet augment G to be a maximal st-graph, call it G^+
- compute upward planar drawings Γ_0^+ and Γ_1^+ of G^+ LR-equivalent to Γ_0 and Γ_1 , respectively

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For any two upward-equivalent planar drawings of an n-vertex plane st-graph G, there exists a O(n)-step upward planar morph between them.

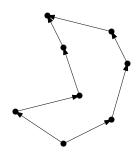
Proof

- ullet augment G to be a maximal st-graph, call it G^+
- compute upward planar drawings Γ_0^+ and Γ_1^+ of G^+ LR-equivalent to Γ_0 and Γ_1 , respectively
- contract an internal low-degree vertex v to one of its neighbors in Γ_0^+ and Γ_1^+ , recursively compute a morph of the obtained drawings, add v back to the obtained morph

Upward Planar Morphs

Reduced Plane st-Graphs Augmentation to G^+ and computation of Γ_0^+ and Γ_1^+

For each internal face f of G:

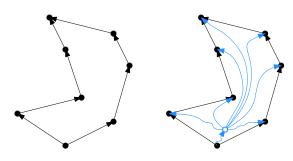


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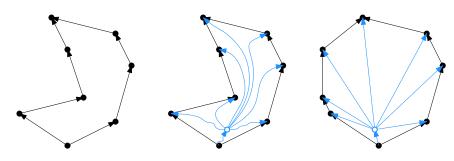
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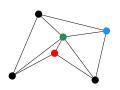


Use the algorithm for drawing hierarchical plane graphs with assigned y-coordinates by Hong & Nagamochi

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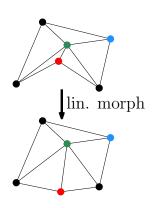
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Edge contractions are widely used in morphing algorithms



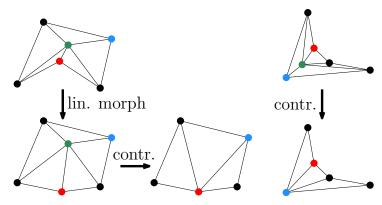


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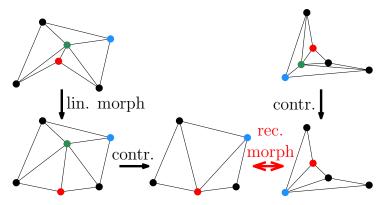
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Edge contractions are widely used in morphing algorithms



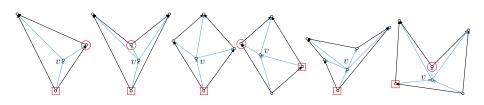
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Reduced Plane st-Graphs Distinguished Neighbors & Contractions

A neighbor u of v is a *distinguished* neighbor of v if it is:

- ullet either a predecessor of v s.t. G has a path from any other predecessor w to v through u, or
- \bullet a successor of v s.t. G has a path from v through u to any other successor w

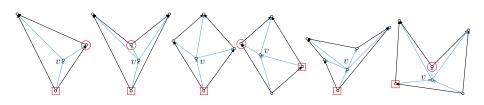


Upward Planar Morphs

Reduced Plane st-Graphs Distinguished Neighbors & Contractions

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With a slight variation of the standard technique, we contract \boldsymbol{v} on a distinguished neighbor

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High-level Idea

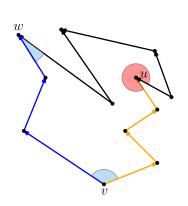
- augment G, Γ_0 , and Γ_1 to obtain a (reduced) st-graph G^+
- ullet compute an upward planar morph ${\mathcal M}$ between Γ_0^+ and Γ_1^+

Theorem

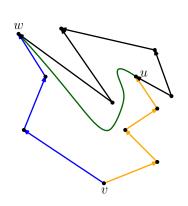
For any two upward equivalent upward planar drawings of an n-vertex upward plane graph G, there exists a $O(n^2)$ -step upward planar morph between them.

High-level Idea

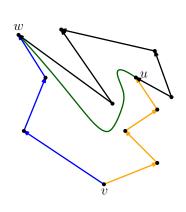
- augment G, Γ_0 , and Γ_1 to obtain a (reduced) st-graph G^+
- \bullet compute an upward planar morph ${\mathcal M}$ between Γ_0^+ and Γ_1^+
- ullet restrict ${\mathcal M}$ to the vertices and edges of G



- G is not an st-graph \implies some faces have *more than one* source and one sink
- Large angles must occur at some of such vertices

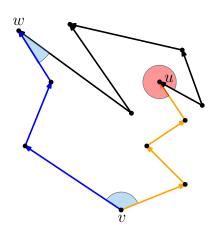


- G is not an st-graph \implies some faces have $more\ than\ one$ source and one sink
- Large angles must occur at some of such vertices
- *splitting* the face at a Large vertex solves the problem locally
- so, add edge <u>uw</u>



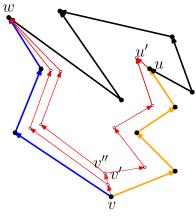
- G is not an st-graph \implies some faces have $more\ than\ one$ source and one sink
- Large angles must occur at some of such vertices
- *splitting* the face at a Large vertex solves the problem locally
- ullet so, add edge uw
- \bullet assume that modifying the drawing for inserting uw costs $T_{uw}(n)$ morphing steps
- ullet G contains O(n) Large vertices
 - \implies turning G into a (reduced) st-graph takes $O(n \cdot T_{uw}(n))$ steps

Removing Large angles



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Removing Large angles



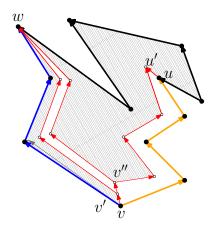
insert the gadget

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Removing Large angles

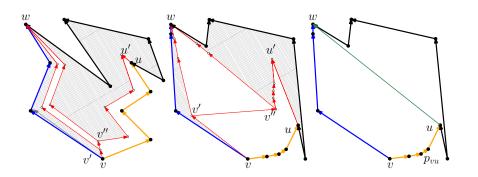


triangulate all the faces except those induced by the gadget

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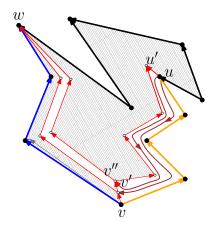
Removing Large angles



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Removing Large angles

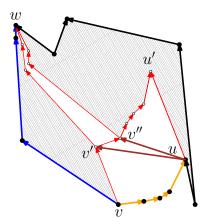


insert edges uv' and uv'' to obtain an st-graph G^* where u is a predecessor of both v' and v''

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Removing Large angles

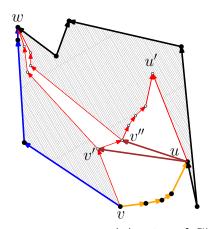


compute an upward drawing of G^*

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Removing Large angles



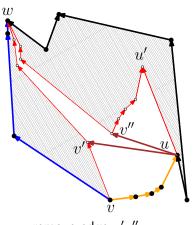
compute an upward drawing of G^{\ast} ...but even by removing $v'v'',\ u$ might not have visibility on w

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Removing Large angles

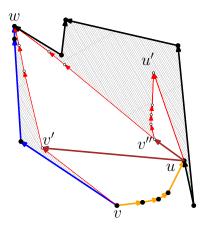


remove edge $v^\prime v^{\prime\prime}$...

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Removing Large angles



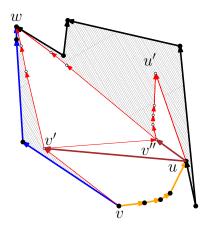
compute a convex upward planar drawing of the obtained st-graph by preserving the y-coordinates (apply Hong & Nagamochi)

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Removing Large angles

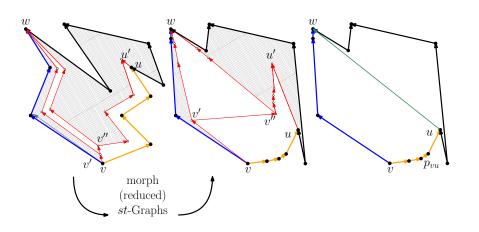


add back $v^\prime v^{\prime\prime}$ and obtain an upward planar drawing of G^*

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Removing Large angles



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Upward Planar Morphs Conclusions and Open Problems

Upward Planar Graphs	$O(n^2)$ – $\Omega(n)$ steps	O(n)?
Reduced Upward Planar Graphs	$\Theta(n)$ steps	
Planar st-Graphs	O(n) steps	O(1)?
Reduced Planar st-Graphs	$\Theta(1)$ steps	
??-Graphs	o(n) steps	

Upward Planar Morphs Conclusions and Open Problems

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Thank you!